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# SSL and IPSec

CS461/ECE422

Fall 2010

# Reading

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- Chapter 11 in Computer Science: Art and Science
- Stallings book
- Can also look at Standards Documents

# SSL

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- Transport layer security
  - Provides confidentiality, integrity, authentication of endpoints
  - Developed by Netscape for WWW browsers and servers
- Internet protocol version: TLS
  - Compatible with SSL
  - Standard [rfc2712](#)

# Working at Transport Level

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- Data link, Network, and Transport headers sent unchanged
- Original transport header can be protected if tunneling

Ethernet Frame Header	IP Header	TCP Header	TCP data stream Encrypted/authenticated Regardless of application
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# SSL Session

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- Association between two peers
  - May have many associated connections
  - Information for each association:
    - Unique session identifier
    - Peer's X.509v3 certificate, if needed
    - Compression method
    - Cipher spec for cipher and MAC
    - “Master secret” shared with peer
      - 48 bits

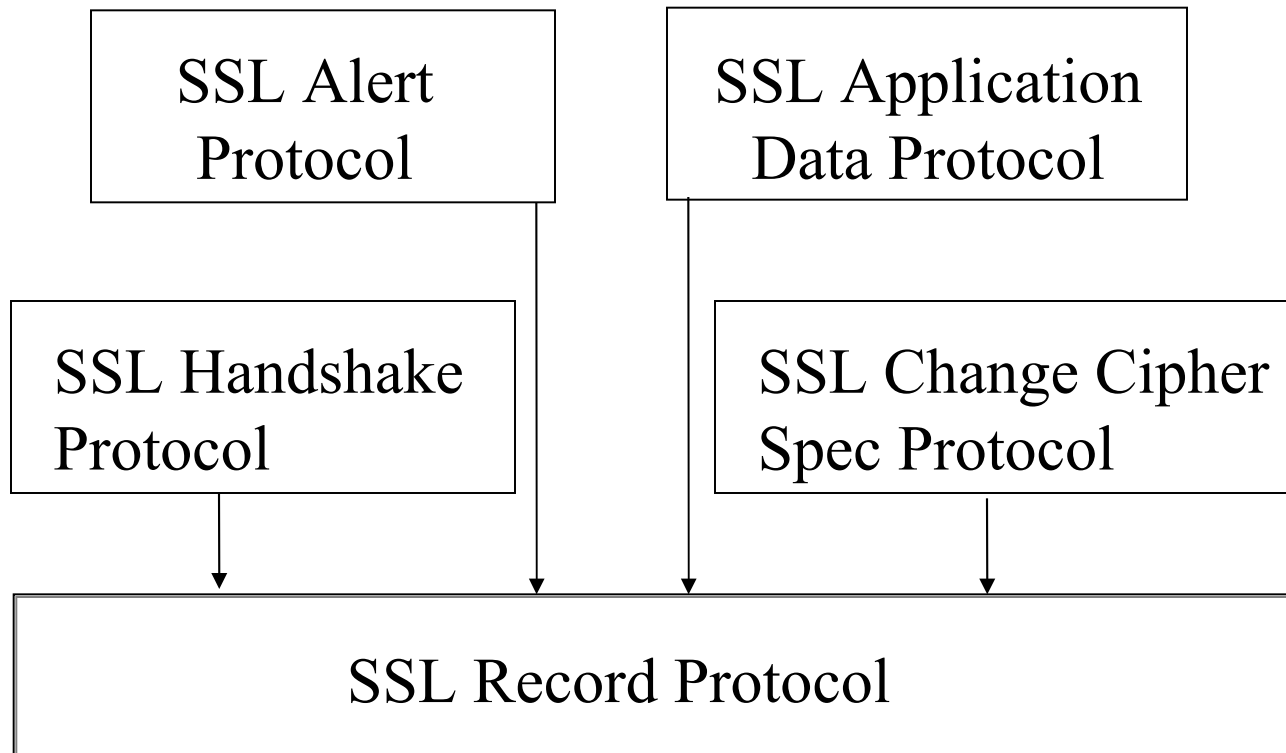
# SSL Connection

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- Describes how data exchanged with peer
- Information for each connection
  - Random data
  - Write keys (used to encipher data)
  - Write MAC key (used to compute MAC)
  - Initialization vectors for ciphers, if needed
  - Sequence numbers

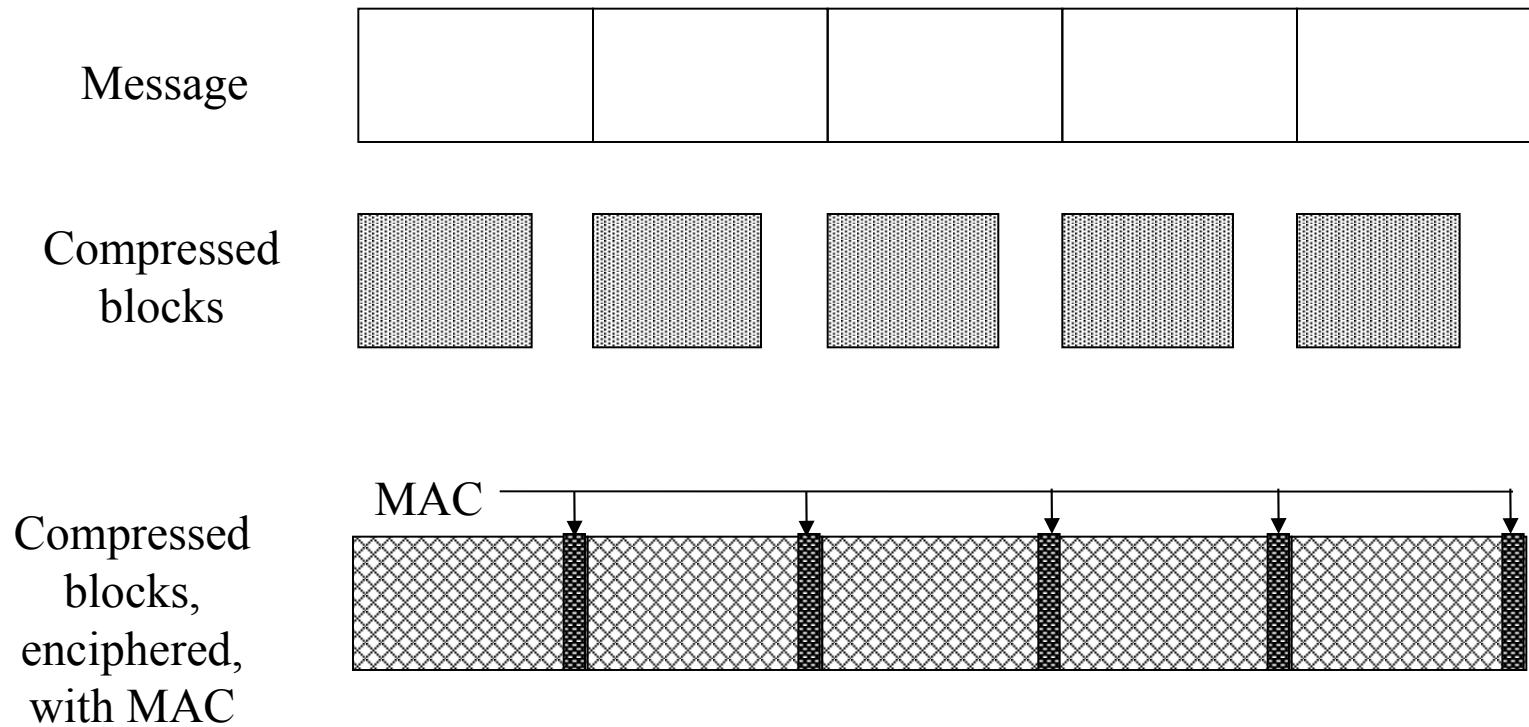
# Structure of SSL

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# SSL Record Layer

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# Record Protocol Overview

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- Lowest layer, taking messages from higher
  - Max block size 16,384 bytes
  - Bigger messages split into multiple blocks
- Construction
  - Block  $b$  compressed; call it  $b_c$
  - MAC computed for  $b_c$ 
    - If MAC key not selected, no MAC computed
  - $b_c$ , MAC enciphered
    - If enciphering key not selected, no enciphering done
  - SSL record header prepended

# SSL MAC Computation

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- Symbols

- $h$  hash function (MD5 or SHA)
- $k_w$  write MAC key of entity
- $ipad = 0x36$ ,  $opad = 0x5C$ 
  - Repeated to block length (from HMAC)
- $seq$  sequence number
- $SSL\_comp$  message type
- $SSL\_len$  block length

- MAC

$h(k_w || opad || h(k_w || ipad || seq || SSL\_comp || SSL\_len || block))$

# SSL Handshake Protocol

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- Used to initiate connection
  - Sets up parameters for record protocol
  - 4 rounds
- Upper layer protocol
  - Invokes Record Protocol
- Note: what follows assumes client, server using RSA as interchange cryptosystem

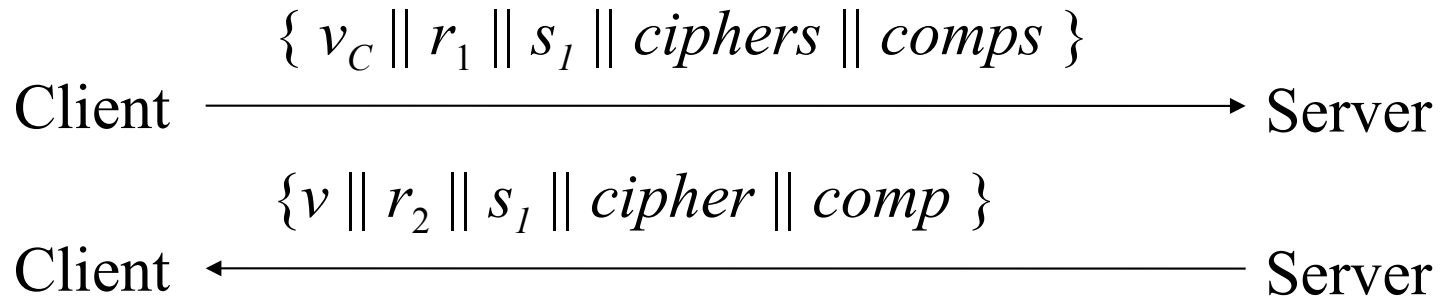
# Overview of Rounds

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- Create SSL connection between client, server
- Server authenticates itself
- Client validates server, begins key exchange
- Acknowledgments all around

# Handshake Round 1

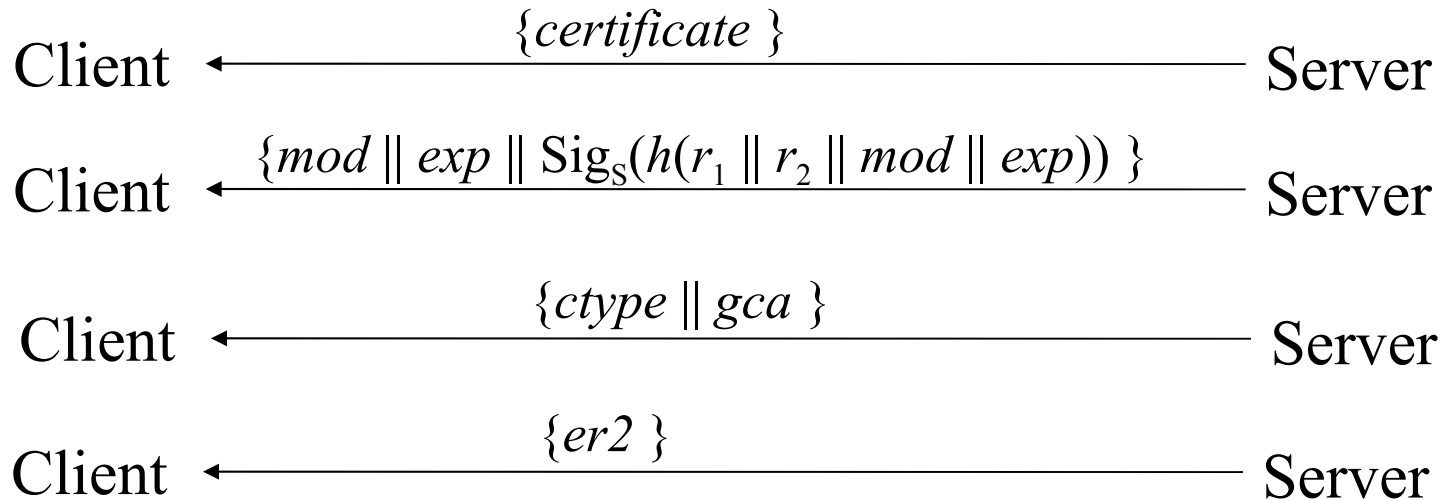
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$v_C$	Client's version of SSL
$v$	Highest version of SSL that Client, Server both understand
$r_1, r_2$	nonces (timestamp and 28 random bytes)
$s_1$	Current session id (0 if new session)
$ciphers$	Ciphers that client understands
$comps$	Compression algorithms that client understand
$cipher$	Cipher to be used
$comp$	Compression algorithm to be used

# Handshake Round 2

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Note: if Server not to authenticate itself, only last message sent; third step omitted if Server does not need Client certificate

$k_s$  Server's private key

$ctype$  Certificate type requested (by cryptosystem)

$gca$  Acceptable certification authorities

$er2$  End round 2 message

# Handshake Round 3

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Client  $\xrightarrow{\{ client\_cert \}}$  Server

Client  $\xrightarrow{\{ pre \} Pub_s}$  Server

Both Client, Server compute master secret *master*:

$master = MD5(pre \parallel SHA('A' \parallel pre \parallel r_1 \parallel r_2) \parallel$

$MD5(pre \parallel SHA('BB' \parallel pre \parallel r_1 \parallel r_2) \parallel$

$MD5(pre \parallel SHA('CCC' \parallel pre \parallel r_1 \parallel r_2))$

Client  $\xrightarrow{\{ h(master \parallel opad \parallel h(msgs \parallel master \parallel ipad)) \}}$  Server

*msgs* Concatenation of previous messages sent/received this handshake

*opad, ipad* As above

# Handshake Round 4

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Client sends “change cipher spec” message using that protocol

Client  $\longrightarrow$  Server

$\{ h(master || opad || h(msgs || 0x434C4E54 || master || ipad)) \}$

Client  $\longrightarrow$  Server

Server sends “change cipher spec” message using that protocol

Client  $\longleftarrow$  Server

$\{ h(master || opad || h(msgs || 0x53525652 || master || ipad)) \}$

Client  $\longleftarrow$  Server

*msgs* Concatenation of messages sent/received this handshake in  
*previous* rounds (does not include these messages)

*opad, ipad, master* As above



# Supporting Crypto

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- All parts of SSL use them
- Initial phase: public key system exchanges keys
  - Classical ciphers ensure confidentiality, cryptographic checksums added for integrity
  - Only certain combinations allowed
    - Depends on algorithm for interchange cipher
  - Interchange algorithms: RSA, Diffie-Hellman, Fortezza
  - AES added in 2002 by [rfc3268](#)

# RSA: Cipher, MAC Algorithms

<i>Interchange cipher</i>	<i>Classical cipher</i>	<i>MAC Algorithm</i>
RSA, key $\leq 512$ bits	<i>none</i>	MD5, SHA
	RC4, 40-bit key	MD5
	RC2, 40-bit key, CBC mode	MD5
	DES, 40-bit key, CBC mode	SHA
RSA	<i>None</i>	MD5, SHA
	RC4, 128-bit key	MD5, SHA
	IDEA, CBC mode	SHA
	DES, CBC mode	SHA
	DES, EDE mode, CBC mode	SHA

# Diffie-Hellman: Types

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- Diffie-Hellman: certificate contains D-H parameters, signed by a CA
  - DSS or RSA algorithms used to sign
- Ephemeral Diffie-Hellman: DSS or RSA certificate used to sign D-H parameters
  - Parameters not reused, so not in certificate
- Anonymous Diffie-Hellman: D-H with neither party authenticated
  - Use is “strongly discouraged” as it is vulnerable to attacks

# D-H: Cipher, MAC Algorithms

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<i>Interchange cipher</i>	<i>Classical cipher</i>	<i>MAC Algorithm</i>
Diffie-Hellman, DSS Certificate	DES, 40-bit key, CBC mode	SHA
	DES, CBC mode	SHA
	DES, EDE mode, CBC mode	SHA
Diffie-Hellman, key $\leq 512$ bits RSA Certificate	DES, 40-bit key, CBC mode	SHA
	DES, CBC mode	SHA
	DES, EDE mode, CBC mode	SHA

# Ephemeral D-H: Cipher, MAC Algorithms

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<i>Interchange cipher</i>	<i>Classical cipher</i>	<i>MAC Algorithm</i>
Ephemeral Diffie-Hellman, DSS Certificate	DES, 40-bit key, CBC mode	SHA
	DES, CBC mode	SHA
	DES, EDE mode, CBC mode	SHA
Ephemeral Diffie-Hellman, key $\leq 512$ bits, RSA Certificate	DES, 40-bit key, CBC mode	SHA
	DES, CBC mode	SHA
	DES, EDE mode, CBC mode	SHA

# Anonymous D-H: Cipher, MAC Algorithms

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<i>Interchange cipher</i>	<i>Classical cipher</i>	<i>MAC Algorithm</i>
Anonymous D-H, DSS Certificate	RC4, 40-bit key	MD5
	RC4, 128-bit key	MD5
	DES, 40-bit key, CBC mode	SHA
	DES, CBC mode	SHA
	DES, EDE mode, CBC mode	SHA

# Fortezza: Cipher, MAC Algorithms

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<i>Interchange cipher</i>	<i>Classical cipher</i>	<i>MAC Algorithm</i>
Fortezza key exchange	<i>none</i>	SHA
	RC4, 128-bit key	MD5
	Fortezza, CBC mode	SHA

# SSL Change Cipher Spec Protocol

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- Send single byte
- In handshake, new parameters considered “pending” until this byte received
  - Old parameters in use, so cannot just switch to new ones



# SSL Alert Protocol

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- Closure alert
  - Sender will send no more messages
  - Pending data delivered; new messages ignored
- Error alerts
  - Warning: connection remains open
  - Fatal error: connection torn down as soon as sent or received

# SSL Alert Protocol Errors

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- Always fatal errors:
  - unexpected\_message, bad\_record\_mac, decompression\_failure, handshake\_failure, illegal\_parameter
- May be warnings or fatal errors:
  - no\_certificate, bad\_certificate, unsupported\_certificate, certificate\_revoked, certificate\_expired, certificate\_unknown

# SSL Application Data Protocol

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- Passes data from application to SSL Record Protocol layer

# MITM Attacks

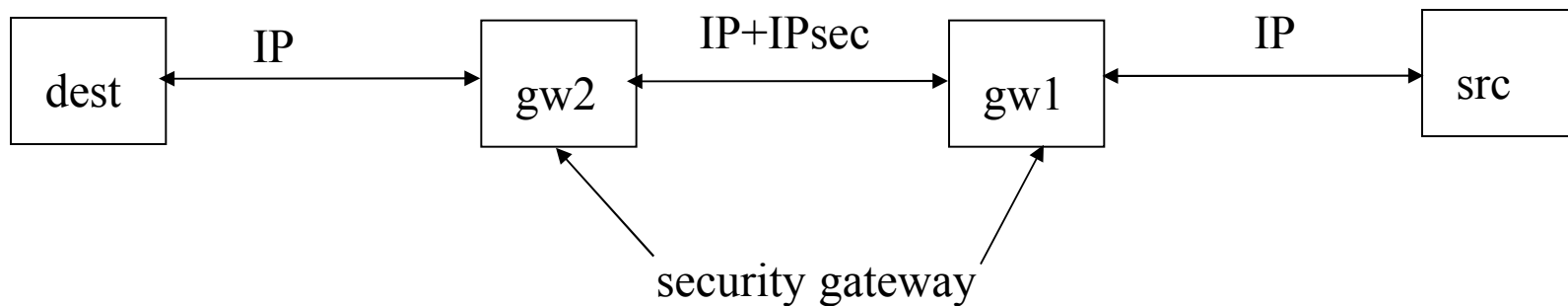
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- Classic attack foiled by certificates
- More subtle attacks appear over time
  - TLS Authentication Gap
    - Interaction of TLS and HTTP
    - <http://www.phonefactor.com/sslgap>
- Application above SSL/TLS tends to be HTTP but does not have to be

# IPsec

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- Network layer security
  - Provides confidentiality, integrity, authentication of endpoints, replay detection
- Protects all messages sent along a path



# Standards

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- Original RFC's 2401-2412
- Mandatory portion of IPv6
- Bolted onto IPv4
- Newer standards
  - IKE: Standardized Key Management Protocol  
RFC 2409
  - NAT-T: UDP encapsulation for traversing  
address translation RFC 3948

# Network Level Encryption

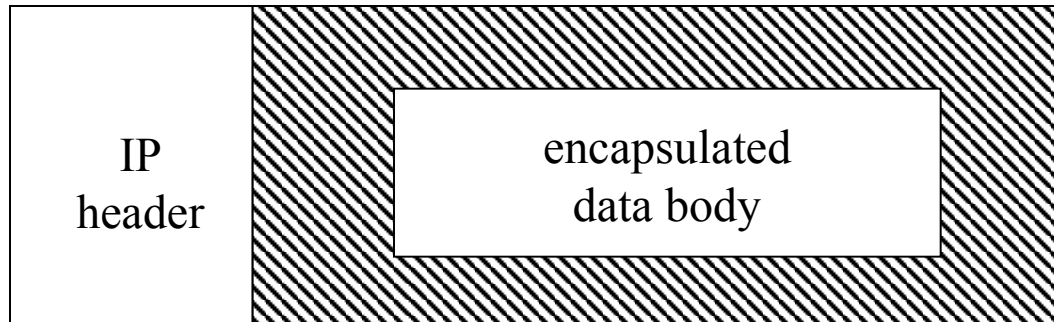
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- Data link header and network header is unchanged
- With tunneling original IP header can be protected

Ethernet Frame Header	IP Header	IP packet Encrypted/authenticated Regardless of application
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# IPsec Transport Mode

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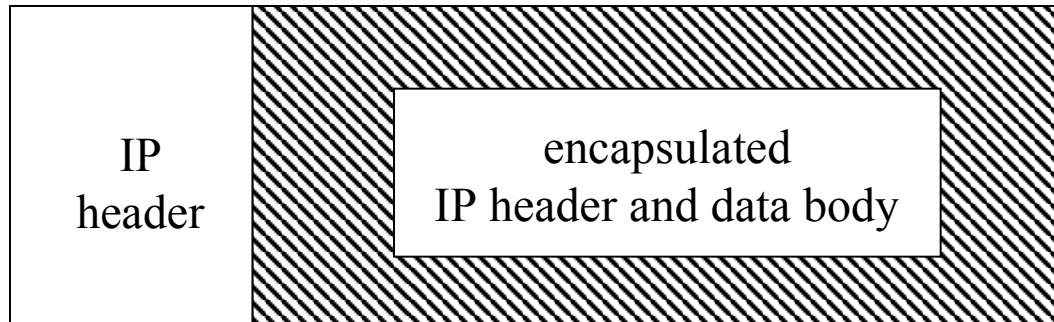


- Encapsulate IP packet data area
- Use IP to send IPsec-wrapped data packet
- Note: IP header not protected



# IPsec Tunnel Mode

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- Encapsulate IP packet (IP header *and* IP data)
- Use IP to send IPsec-wrapped packet
- Note: IP header protected

# IPsec Protocols

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- Authentication Header (AH)
  - Integrity of payload
  - Integrity of outer header
  - Anti-replay
- Encapsulating Security Payload (ESP)
  - Confidentiality of payload and inner header
  - Integrity of payload (and now header)

# ESP and integrity

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- Originally design, use AH to add integrity if needed.
- Bellovin showed integrity is always needed
  - So added directly to ESP
  - <http://www.cs.columbia.edu/~smb/papers/bade>

# IPsec Architecture

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- Security Policy Database (SPD)
  - Says how to handle messages (discard them, add security services, forward message unchanged)
  - SPD associated with network interface
  - SPD determines appropriate entry from packet attributes
    - Including source, destination, transport protocol

# Example

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- Goals
  - Discard SMTP packets from host 192.168.2.9
  - Forward packets from 192.168.19.7 without change

- SPD entries

```
src 192.168.2.9, dest 10.1.2.3 to 10.1.2.103, port 25, discard
src 192.168.19.7, dest 10.1.2.3 to 10.1.2.103, port 25, bypass
dest 10.1.2.3 to 10.1.2.103, port 25, apply IPsec
```

- Note: entries scanned in order
  - If no match for packet, it is discarded

# IPsec Architecture

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- Security Association (SA)
  - Association between peers for security services
    - Identified uniquely by dest address, security protocol (AH or ESP), unique 32-bit number (security parameter index, or SPI)
  - Unidirectional
    - Can apply different services in either direction
  - SA uses either ESP or AH; if both required, 2 SAs needed

# SA Database (SAD)

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- Entry describes SA; some fields for all packets:
  - AH algorithm identifier, keys
    - When SA uses AH
  - ESP encipherment algorithm identifier, keys
    - When SA uses confidentiality from ESP
  - ESP authentication algorithm identifier, keys
    - When SA uses authentication, integrity from ESP
  - SA lifetime (time for deletion or max byte count)
  - IPsec mode (tunnel, transport, either)

# SAD Fields

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- Antireplay (inbound only)
  - When SA uses antireplay feature
- Sequence number counter (outbound only)
  - Generates AH or ESP sequence number
- Sequence counter overflow field
  - Stops traffic over this SA if sequence counter overflows
- Aging variables
  - Used to detect time-outs



# IPsec Architecture

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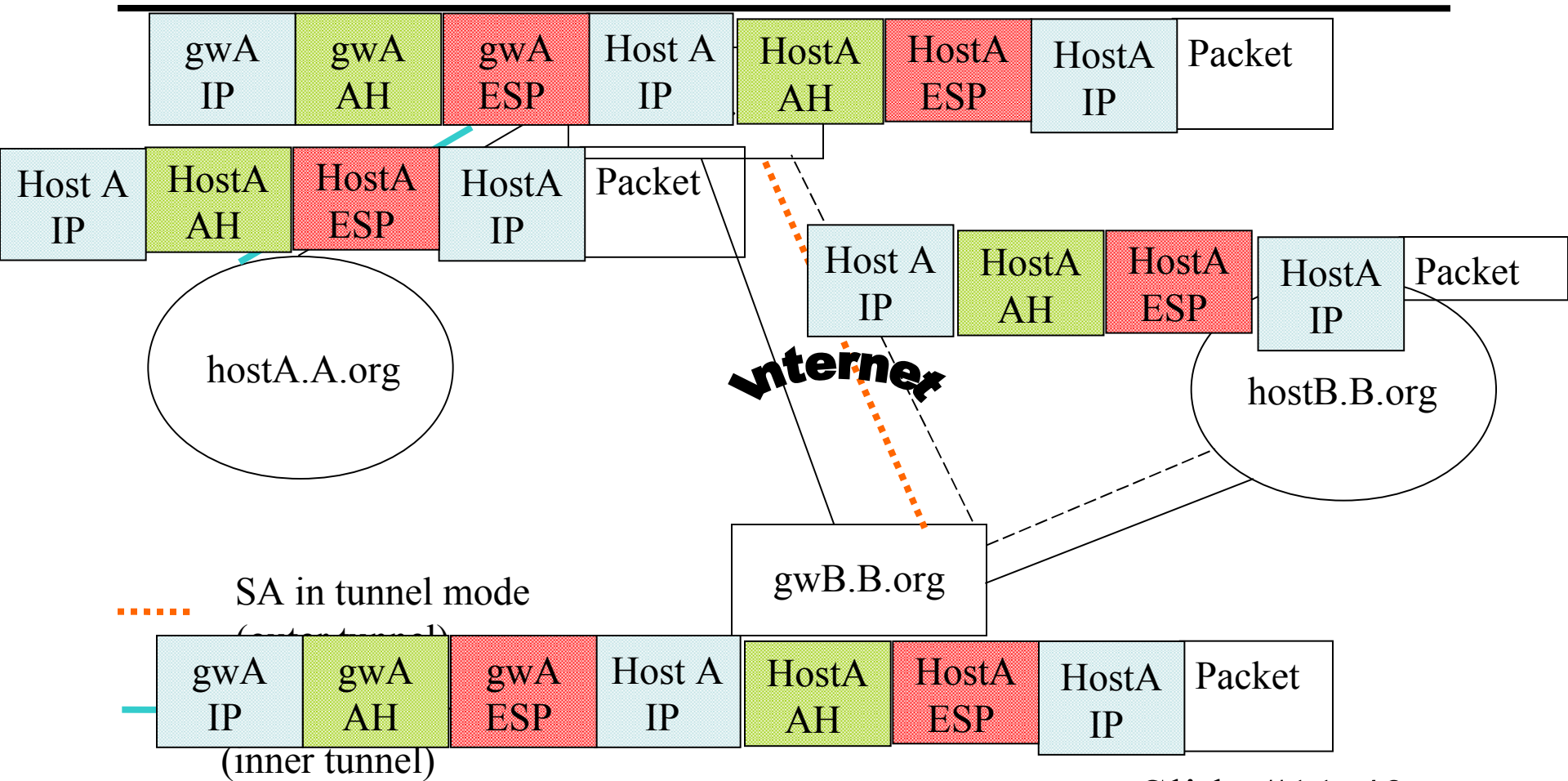
- Packet arrives
- Look for existing SA
- Otherwise look in SPD
  - Find appropriate entry
  - Get dest address, security protocol, SPI
- Find associated SA in SAD
  - Use dest address, security protocol, SPI
  - Apply security services in SA (if any)

# Example: Nested Tunnels

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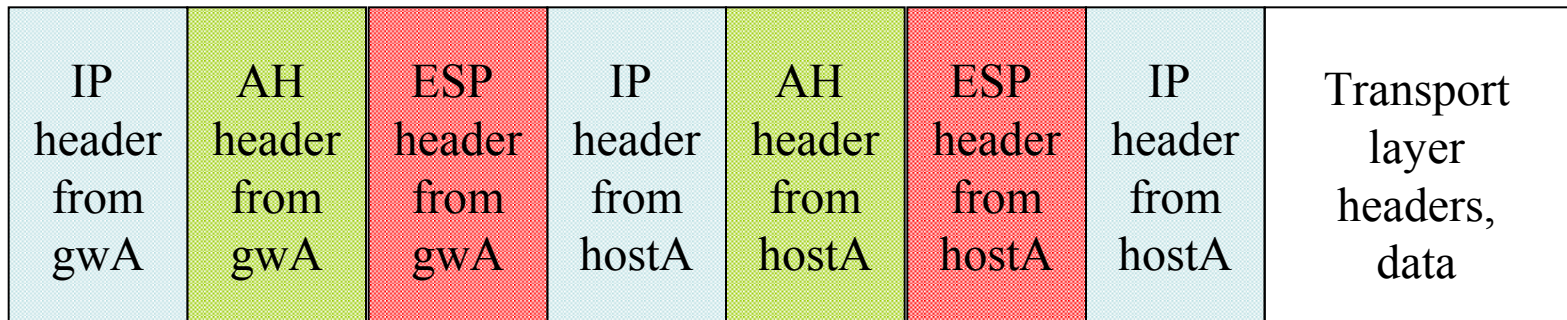
- Group in A.org needs to communicate with group in B.org
- Gateways of A, B use IPsec mechanisms
  - But the information must be secret to everyone except the two groups, even secret from other people in A.org and B.org
- Inner tunnel: a SA between the hosts of the two groups
- Outer tunnel: the SA between the two gateways

# Example: Systems



# Example: Packets

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- Packet generated on hostA
- Encapsulated by hostA's IPsec mechanisms
- Again encapsulated by gwA's IPsec mechanisms
  - Above diagram shows headers, but as you go left, everything to the right would be enciphered and authenticated, *etc.*

# AH Protocol

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- Parameters in AH header
  - Length of header
  - SPI of SA applying protocol
  - Sequence number (anti-replay)
  - Integrity value check
- Two steps
  - Check that replay is not occurring
  - Check authentication data

# Sender

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- Check sequence number will not cycle
- Increment sequence number
- Compute IVC of packet
  - Includes IP header, AH header, packet data
    - IP header: include all fields that will not change in transit; assume all others are 0
    - AH header: authentication data field set to 0 for this
    - Packet data includes encapsulated data, higher level protocol data

# Recipient

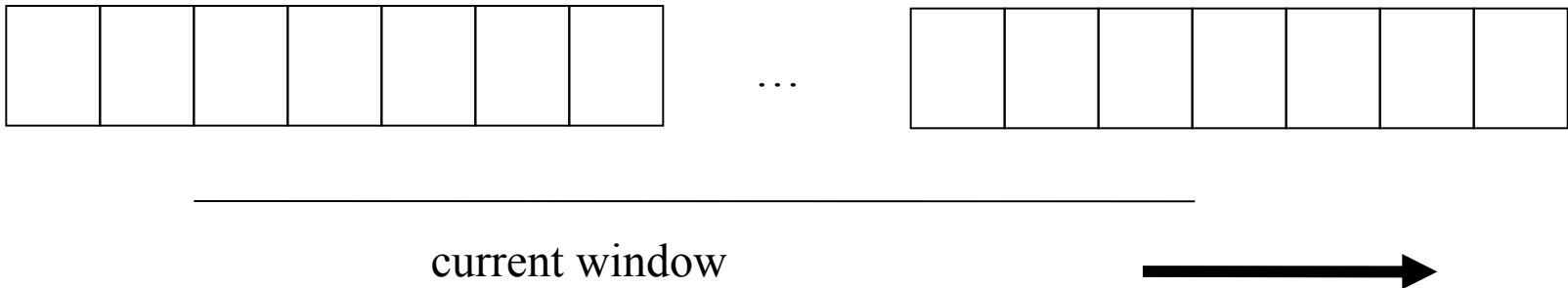
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- Assume AH header found
- Get SPI, destination address
- Find associated SA in SAD
  - If no associated SA, discard packet
- If antireplay not used
  - Verify IVC is correct
    - If not, discard

# Recipient, Using Antireplay

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- Check packet beyond low end of sliding window
- Check IVC of packet
- Check packet's slot not occupied
  - If any of these is false, discard packet





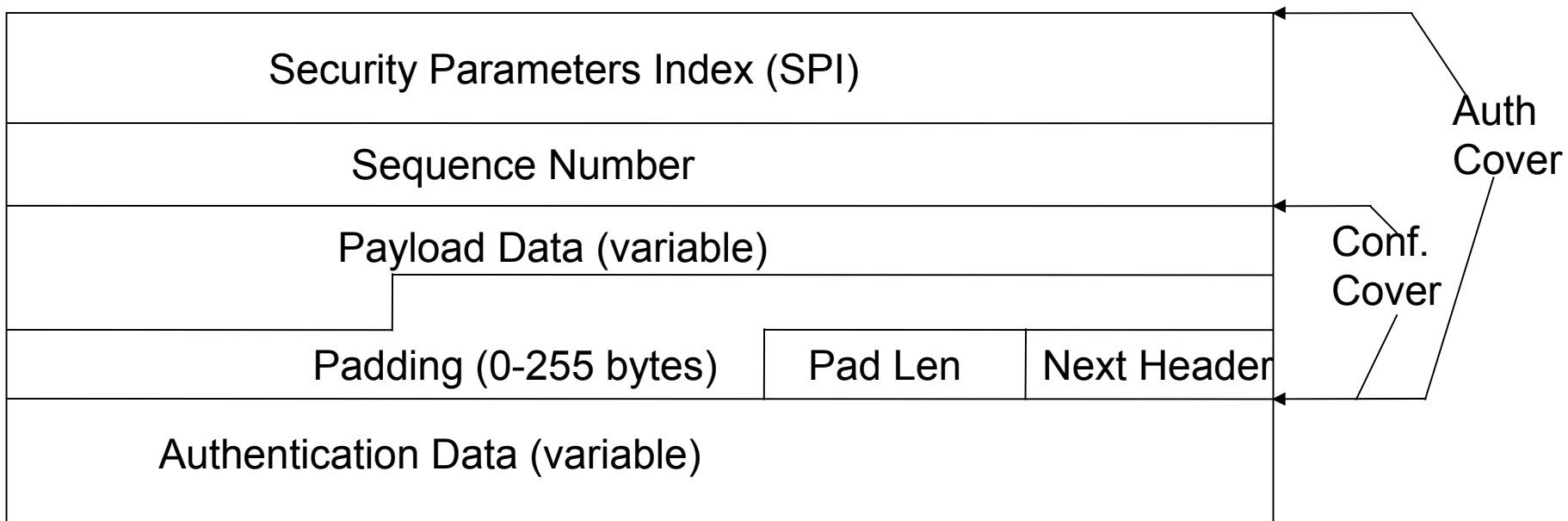
# AH Miscellany

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- All implementations must support:  
HMAC\_MD5  
HMAC\_SHA-1
- May support other algorithms

# ESP Header

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# ESP Protocol

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- Parameters in ESP header
  - SPI of SA applying protocol
  - Sequence number (anti-replay)
  - Generic “payload data” field
  - Padding and length of padding
    - Contents depends on ESP services enabled; may be an initialization vector for a chaining cipher, for example
    - Used also to pad packet to length required by cipher
  - Optional authentication data field

# Sender

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- Add ESP header
  - Includes whatever padding needed
- Encipher result
  - Do not encipher SPI, sequence numbers
- If authentication desired, compute as for AH protocol *except* over ESP header, payload and *not* encapsulating IP header

# Recipient

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- Assume ESP header found
- Get SPI, destination address
- Find associated SA in SAD
  - If no associated SA, discard packet
- If authentication used
  - Do IVC, antireplay verification as for AH
    - Only ESP, payload are considered; *not* IP header
    - Note authentication data inserted after encipherment, so no deciphering need be done

# Recipient

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- If confidentiality used
  - Decipher enciphered portion of ESP header
  - Process padding
  - Decipher payload
  - If SA is transport mode, IP header and payload treated as original IP packet
  - If SA is tunnel mode, payload is an encapsulated IP packet and so is treated as original IP packet

# ESP Miscellany

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- Must use at least one of confidentiality, authentication services
- Synchronization material must be in payload
  - Packets may not arrive in order, so if not, packets following a missing packet may not be decipherable

# More ESP Miscellany

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- All implementations must support (encipherment algorithms):

DES in CBC mode

NULL algorithm (identity; no encipherment)

- All implementations must support (integrity algorithms):

HMAC\_MD5

HMAC\_SHA-1

NULL algorithm (no MAC computed)

- Both cannot be NULL at the same time



# Which to Use: PEM, SSL, IPsec

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- What do the security services apply to?
  - If applicable to one application *and* application layer mechanisms available, use that
    - PEM for electronic mail
  - If more generic services needed, look to lower layers
    - SSL for transport layer, end-to-end mechanism
    - IPsec for network layer, either end-to-end or link mechanisms, for connectionless channels as well as connections
  - If endpoint is host, SSL and IPsec sufficient; if endpoint is user, application layer mechanism such as PEM needed

# Key Points

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- Key management critical to effective use of cryptosystems
  - Different levels of keys (session vs. interchange)
- Keys need infrastructure to identify holders, allow revoking
  - Key escrowing complicates infrastructure
- Digital signatures provide integrity of origin and content

Much easier with public key cryptosystems than with classical cryptosystems