

# Programming Languages and Compilers (CS 421)



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<http://courses.engr.illinois.edu/cs421/su2013/>

Based in part on slides by Mattox Beckman, as updated by Vikram Adve, Gul Agha, and Elsa Gunter



# Recursive Functions

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```
# let rec factorial n =  
    if n = 0 then 1 else n * factorial (n - 1);;  
val factorial : int -> int = <fun>  
# factorial 5;;  
- : int = 120  
# (* rec is needed for recursive function  
   declarations *)
```



# Recursion Example

Compute  $n^2$  recursively using:

$$n^2 = (2 * n - 1) + (n - 1)^2$$

```
# let rec nthsq n =          (* rec for recursion *)
  match n                    (* pattern matching for cases *)
  with 0 -> 0                (* base case *)
  | n -> (2 * n - 1)          (* recursive case *)
      + nthsq (n - 1);;      (* recursive call *)
val nthsq : int -> int = <fun>
# nthsq 3;;
- : int = 9
```

Structure of recursion similar to inductive proof



# Recursion and Induction

---

```
# let rec nthsq n = match n with 0 -> 0  
  | n -> (2 * n - 1) + nthsq (n - 1) ;;
```

- Base case is the last case; it stops the computation
- Recursive call must be to arguments that are somehow smaller - must progress to base case
- **if** or **match** must contain base case
- Failure of these may cause failure of termination

- Simple recursive (“algebraic”) datatype
- Unlike tuples, lists are type homogeneous (all elements same type) but have varying length

- List can take one of two forms:
  - Empty list, written `[ ]`
  - Non-empty list, written `x :: xs`
    - `x` is head element, `xs` is tail list, `::` called “cons”
  - Syntactic sugar: `[x] == x :: [ ]`
  - `[ x1; x2; ...; xn ] == x1 :: x2 :: ... :: xn :: [ ]`



# Lists

---

```
# let fib5 = [8;5;3;2;1;1];;
```

```
val fib5 : int list = [8; 5; 3; 2; 1; 1]
```

```
# let fib6 = 13 :: fib5;;
```

```
val fib6 : int list = [13; 8; 5; 3; 2; 1; 1]
```

```
# (8::5::3::2::1::1::[ ]) = fib5;;
```

```
- : bool = true
```

```
# fib5 @ fib6;;
```

```
- : int list = [8; 5; 3; 2; 1; 1; 13; 8; 5; 3; 2; 1; 1]
```



# Lists are Homogeneous

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```
# let bad_list = [1; 3.2; 7];;
```

Characters 19-22:

```
let bad_list = [1; 3.2; 7];;  
                ^^^
```

This expression has type float but is here  
used with type int





# Question

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■ Which one of these lists is invalid?

1. [2; 3; 4; 6]
2. [2,3; 4,5; 6,7]
3. [(2.3, 4); (3.2, 5); (6, 7.2)]
4. [[“hi”; “there”]; [“whatcha”]; [ ]; [“doin”]]



# Answer

---

- Which one of these lists is invalid?

1. [2; 3; 4; 6]
2. [2,3; 4,5; 6,7]
3. [(2.3, 4); (3.2, 5); (6, 7.2)]
4. [[“hi”; “there”]; [“whatcha”]; [ ]; [“doin”]]

- 3 is invalid because of last pair



# Structural Recursion : List Example

```
# let rec length list = match list
  with [ ] -> 0    (* Nil case *)
       | (x :: xs) -> 1 + length xs;; (* Cons case *)
val length : 'a list -> int = <fun>
# length [5; 4; 3; 2];;
- : int = 4
```

- Nil case [ ] is base case
- Cons case recurses on component list xs



# Structural Recursion

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- Functions on recursive datatypes (e.g. lists) tend to be recursive
- Recursion over recursive datatypes generally by structural recursion
  - Recursive calls made to components of structure of the same recursive type
  - Base cases of recursive types stop the recursion of the function



# Functions Over Lists

---

```
# let rec double_up list =  
  match list  
  with [ ] -> [ ] (* pattern before ->,  
                    expression after *)  
       | (x :: xs) -> (x :: x :: double_up xs);;  
val double_up : 'a list -> 'a list = <fun>  
# let fib5_2 = double_up fib5;;  
val fib5_2 : int list = [8; 8; 5; 5; 3; 3; 2; 2; 1;  
  1; 1; 1]
```



# Functions Over Lists

---

```
# let words = double_up ["hi"; "there"];;  
val words : string list = ["hi"; "hi"; "there"; "there"]  
# let rec rev1 list =  
  match list  
  with [] -> []  
       | (x::xs) -> rev1 xs @ [x];; (* add x at the end *)  
val rev1 : 'a list -> 'a list = <fun>  
# rev1 words;;  
- : string list = ["there"; "there"; "hi"; "hi"]
```



# Functions Over Lists

---

```
# let rec map f list =  
  match list  
  with [] -> []  
       | (h::t) -> (f h) :: (map f t);;  
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>  
# map plus_two fib5;;  
- : int list = [10; 7; 5; 4; 3; 3]  
# map (fun x -> x - 1) fib6;;  
- : int list = [12; 7; 4; 2; 1; 0; 0]
```



# Mapping Recursion

- One common form of structural recursion applies a function to each element in the structure

```
# let rec doubleList list = match list
  with [ ] -> [ ]
       | x::xs -> 2 * x :: doubleList xs;;
val doubleList : int list -> int list = <fun>
# doubleList [2;3;4];;
- : int list = [4; 6; 8]
```





# Mapping Recursion

- Can use the higher-order recursive map function instead of direct recursion

```
# let doubleList list =  
  List.map (fun x -> 2 * x) list;;  
val doubleList : int list -> int list = <fun>  
# doubleList [2;3;4];;  
- : int list = [4; 6; 8]
```

- Same function, but no rec



# Iterating over lists

---

```
# let rec fold_left f a list =  
  match list  
  with [] -> a  
       | (x :: xs) -> fold_left f (f a x) xs;;  
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a =  
  <fun>  
# fold_left  
  (fun () -> fun s -> print_string s)  
  ()  
  ["hi"; "there"];;  
hithere- : unit = ()
```



# Iterating over lists

---

```
# let rec fold_right f list b =  
  match list  
  with [] -> b  
       | (x :: xs) -> f x (fold_right f xs b);;  
val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b =  
  <fun>  
# fold_right  
  (fun s -> fun () -> print_string s)  
  ["hi"; "there"]  
  ();;  
therehi- : unit = ()
```



# Folding Recursion

- Another common recursion pattern “folds” an operation over elements of the structure

```
# let rec multList list = match list
  with [ ] -> 1
       | x::xs -> x * multList xs;;
val multList : int list -> int = <fun>
# multList [2;4;6];;
- : int = 48
```

- Computes:  $(2 * (4 * (6 * 1)))$

# Encoding Recursion with Fold

```
# let rec multList list = match list with  
  [ ] -> 1 | x::xs -> x * multList xs;;  
val multList : int list -> int = <fun>
```

Base Case

Operation

Recursive Call

```
# let multList list =  
  fold_right (fun x p -> x * p) list 1;;  
val append : 'a list -> 'a list -> 'a list = <fun>  
# multList [2;4;6];;  
- : int = 48
```



# Forward Recursion

---

- Structural recursion: split input into components, recurse
- One kind of structural recursion is Forward Recursion: recurse at the front
  - Split input into components
  - Recursive call on all recursive components
  - Build final result from partial results
- Wait until the whole structure has been traversed to start building the answer



# Forward Recursion: Examples

---

```
# let rec double_up list =  
  match list  
  with [ ] -> [ ]  
       | (x :: xs) -> (x :: x :: double_up xs);;  
val double_up : 'a list -> 'a list = <fun>
```

```
# let rec rev1 list =  
  match list  
  with [] -> []  
       | (x::xs) -> rev1 xs @ [x];;  
val rev1 : 'a list -> 'a list = <fun>
```



# Forward Recursion: More Examples

---

```
# let rec addList list = match list with  
  [ ] -> 0 | x::xs -> x + addList xs;;
```

```
val addList : int list -> int = <fun>
```

```
# addList [2;3;4];;
```

```
- : int = 9
```

```
# let rec multList list = match list with  
  [ ] -> 1 | x::xs -> x * multList xs;;
```

```
val multList : int list -> int = <fun>
```

```
# multList [2;3;4];;
```

```
- : int = 24
```





# Folding - Forward Recursion

---

```
# let addList list = fold_right (+) list 0;;
```

```
val addList : int list -> int = <fun>
```

```
# addList [2;3;4];;
```

```
- : int = 9
```

```
# let multList list = fold_right ( * ) list 1;;
```

```
val multList : int list -> int = <fun>
```

```
# multList [2;3;4];;
```

```
- : int = 24
```

fold\_right encapsulates forward recursion



# How long will it take?

---

- Recall the big-O notation from CS 225 and CS 273
- Question: given input of size  $n$ , how long to generate output?
- Express output time in terms of input size, omit constants and take biggest power



# How long will it take?

---

Common big-O times:

- Constant time  $O(1)$ 
  - input size doesn't matter
- Linear time  $O(n)$ 
  - double input  $\Rightarrow$  double time
- Quadratic time  $O(n^2)$ 
  - double input  $\Rightarrow$  quadruple time
- Exponential time  $O(2^n)$ 
  - increment input  $\Rightarrow$  double time



# Linear Time

---

- Expect most list operations to take linear time  $O(n)$
- Each step of the recursion can be done in constant time
- Each step makes only one recursive call
- List example: `multList`, `append`
- Integer example: `factorial`



# Quadratic Time

---

- Each step of the recursion takes time proportional to input
- Each step of the recursion makes only one recursive call.
- List example:

```
# let rec rev1 list = match list  
  with [] -> []  
       | (x::xs) -> rev1 xs @ [x];;  
val rev1 : 'a list -> 'a list = <fun>
```



# Exponential running time

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- Hideous running times on input of any size
- Each step of recursion takes constant time
- Each recursion makes two recursive calls
- Easy to accidentally write exponential code for functions that can be linear



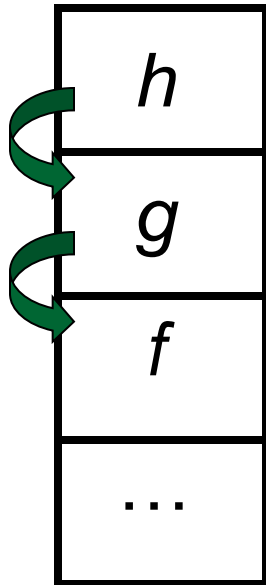
# Exponential running time

---

```
# let rec naiveFib n = match n
  with 0 -> 0
    | 1 -> 1
    | _ -> naiveFib (n-1) + naiveFib (n-2);;
val naiveFib : int -> int = <fun>
```

# Writing Fast Functions

Normal  
call

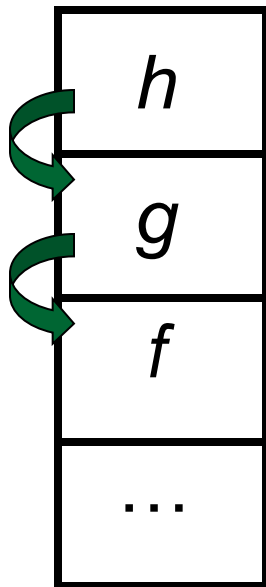


- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished



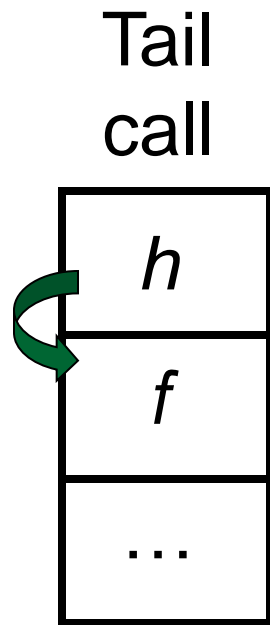
# An Important Optimization

Normal  
call



- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if *f* calls *g* and *g* calls *h*, but calling *h* is the last thing *g* does (a *tail call*)?

# An Important Optimization



- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if *f* calls *g* and *g* calls *h*, but calling *h* is the last thing *g* does (a *tail call*)?
- Then *h* can return directly to *f* instead of *g*



# Tail Recursion

---

- A recursive program is tail recursive if all recursive calls are tail calls
- Tail recursive programs may be implemented as loops, removing the function call overhead for the recursive calls
- Tail recursion generally requires extra “accumulator” arguments to pass partial results – build answer as we go
  - May require an auxiliary function



# Tail Recursion - Example

---

```
# let rec rev_aux list revlist =  
  match list with [ ] -> revlist  
  | x :: xs -> rev_aux xs (x::revlist);;  
val rev_aux : 'a list -> 'a list -> 'a list = <fun>
```

```
# let rev list = rev_aux list [ ];;  
val rev : 'a list -> 'a list = <fun>
```

- What is its running time?



# Tail Recursion - Example

---

```
# let rec rev_aux list revlist =  
  match list with [ ] -> revlist  
  | x :: xs -> rev_aux xs (x::revlist);;  
val rev_aux : 'a list -> 'a list -> 'a list = <fun>
```

```
# let rev list = rev_aux list [ ];;  
val rev : 'a list -> 'a list = <fun>
```

- What is its running time?

$O(n)$



# Comparison

---

- $\text{rev1 } [1,2,3] =$
- $(\text{rev1 } [2,3]) @ [1] =$
- $((\text{rev1 } [3]) @ [2]) @ [1] =$
- $(((\text{rev1 } []) @ [3]) @ [2]) @ [1] =$
- $(([] @ [3]) @ [2]) @ [1] =$
- $([3] @ [2]) @ [1] =$
- $(3 :: ([ ] @ [2])) @ [1] =$
- $[3,2] @ [1] =$
- $3 :: ([2] @ [1]) =$
- $3 :: (2 :: ([ ] @ [1])) = [3, 2, 1]$



# Comparison

---

- $\text{rev}[1,2,3] =$
- $\text{rev\_aux}[1,2,3][ ] =$
- $\text{rev\_aux}[2,3][1] =$
- $\text{rev\_aux}[3][2,1] =$
- $\text{rev\_aux}[ ][3,2,1] = [3,2,1]$



# Folding

---

```
# let rec fold_left f a list = match list
  with [] -> a | (x :: xs) -> fold_left f (f a x) xs;;
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a =
  <fun>
```

```
fold_left f a [x1; x2;...;xn] = f (...(f (f a x1) x2)...) xn
```

```
# let rec fold_right f list b = match list
  with [ ] -> b | (x :: xs) -> f x (fold_right f xs b);;
val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b =
  <fun>
```

```
fold_right f [x1; x2;...;xn] b = f x1 (f x2 (...(f xn b)...))
```





# Folding - Tail Recursion

---

```
let rec rev_aux list revlist =  
  match list with [ ] -> revlist  
  | x :: xs -> rev_aux xs (x::revlist);;  
let rev list = rev_aux list [ ];;
```

```
# let rev list =  
  fold_left  
    (fun l x -> x :: l)    //comb op  
    []                    //accumulator cell  
  list;;
```



# Folding

---

- Can replace recursion by `fold_right` in most forward recursive definitions
- Can replace recursion by `fold_left` in most tail recursive definitions



# Map from Fold

---

```
# let map f list =  
  fold_right (fun x y -> f x :: y) list [ ];;  
val map : ('a -> 'b) -> 'a list -> 'b list =  
  <fun>  
  
# map ((+) 1) [1;2;3];;  
- : int list = [2; 3; 4]
```

- Can you write fold\_right (or fold\_left) with just map? How, or why not?

# Higher Order Functions

- A function is *higher-order* if it takes a function as an argument or returns one as a result
- Example:

```
# let compose f g = fun x -> f (g x);;
```

```
val compose : ('a -> 'b) -> ('c -> 'a) -> 'c -> 'b = <fun>
```

- The type  $('a \rightarrow 'b) \rightarrow ('c \rightarrow 'a) \rightarrow 'c \rightarrow 'b$  is a higher order type because of  $('a \rightarrow 'b)$  and  $('c \rightarrow 'a)$  and  $\rightarrow 'c \rightarrow 'b$



# Partial Application

---

```
# (+);;
```

```
- : int -> int -> int = <fun>
```

```
# (+) 2 3;;
```

```
- : int = 5
```

```
# let plus_two = (+) 2;;
```

```
val plus_two : int -> int = <fun>
```

```
# plus_two 7;;
```

```
- : int = 9
```

■ Patial application also called *sectioning*



# Lambda Lifting

---

- You must remember the rules for evaluation when you use partial application

```
# let add_two = (+) (print_string "test\n"; 2);;
```

```
test
```

```
val add_two : int -> int = <fun>
```

```
# let add2 = (* lambda lifted *)
```

```
  fun x -> (+) (print_string "test\n"; 2) x;;
```

```
val add2 : int -> int = <fun>
```



# Lambda Lifting

---

```
# thrice add_two 5;;
```

```
- : int = 11
```

```
# thrice add2 5;;
```

```
test
```

```
test
```

```
test
```

```
- : int = 11
```

- Lambda lifting delayed the evaluation of the argument to (+) until the second argument was supplied



# Partial Application and “Unknown Types”

- Consider compose plus\_two:

```
# let f1 = compose plus_two;;
```

```
val f1 : ('_a -> int) -> '_a -> int = <fun>
```

- Compare to lambda lifted version:

```
# let f2 = fun g -> compose plus_two g;;
```

```
val f2 : ('a -> int) -> 'a -> int = <fun>
```

- What is the difference?



# Partial Application and “Unknown Types”

- ‘\_a can only be instantiated once for an expression

```
# f1 plus_two;;
```

```
- : int -> int = <fun>
```

```
# f1 List.length;;
```

Characters 3-14:

```
f1 List.length;;
```

```
^^^^^^^^^^
```

This expression has type 'a list -> int but is here used  
with type int -> int



# Partial Application and “Unknown Types”

---

- ‘a can be repeatedly instantiated

```
# f2 plus_two;;
```

```
- : int -> int = <fun>
```

```
# f2 List.length;;
```

```
- : 'a list -> int = <fun>
```