CS 421 — Programming Language and Compilers

Welcome!

In today's class:

- Course intro
 - What CS 421 is about
 - Class structure
- Brief discussion of programming language history
- Intro to the OCaml programming language

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What you will learn this semester

- How to implement programming languages
 - Writing lexical analyzers and parsers
 - Translating programs to machine language
 - Implementing run-time systems
- How to write programs in a functional programm language
- How to formally define languages (including the definiti of type rules and of program execution)
- Key differences between statically-typed languages (e.g. Java) and dynamically-typed languages (Python, JavaScr.
- Plus a few other things...

Why learn about lexing and parsing?

Learn basic algorithms for dealing with structus textual data, applying theory from CS 373:

- DFAs and regular expressions for "lexing"
- Top-down (aka recursive descent) parsing
- Bottom-up (aka LR, aka shift/reduce) parsing

Why learn about definition and implementation of conventional languages?

- Understand standard languages more precisely
- Complete picture of how programs go from keyboard execution
- Understand translation from high-level language machine language — help understand inefficiencies
- Learn to build language processors (esp. compilers)
- Learn basic concepts for handling structured text day
 abstract syntax; syntax-directed translation

Why learn about definition, implementation, and use of functional languages?

- Functional programming is the major alternative program paradigm to imperative/object-oriented programming.
- Functional programming language features are increasir appearing in mainstream languages
- Functional programming concepts increasingly used in n functional languages, and especially in scripting languages easiest to learn them in a functional programming languages

How this class operates

- Lectures 9:30–10:45 TR in 1404 SC
- Usually weekly assignments, due Tuesday at 9:30am
- Programming mostly in OCaml, a functional language
- Each class's slides put on web day before class; subsethe slides the ones containing exercises printed a distributed in class
- Two midterms and final (dates t.b.a.)
- Course policies grade calculations, collaboration/chea policy, lateness — given on web

Classroom pedagogy

- In-class exercises (ungraded) These are smaller version what you'll be doing on homeworks and exams.
- Discussion encouraged when doing exercises
- Tablet PCs will be used to help me see what you're doing possibly share with class

Brief overview of programming languages

- First high-level language (FORTRAN) developed c. 19
 Emphasis was on efficiency.
- As computers got faster, and memories grew, langua became less and less about efficiency, more and more ab programmer productivity.
- Modern language design involves trade-offs among vari issues: efficiency; (short-term) programmer convenier (long-term) program maintainability; portability; secur parallelizability.

Capsule history of PLs

	Conventional	О-о	Scripting	Functional
1957/8	Fortran	<u> </u>	beripung	LISP
1960's	Algol	Simula 67		
19 7 0's	C	Smalltalk		ML
1980's		C++, Obj. C		
1990's		Java	Perl	Haskell
			JavaScript	OCaml
early 2000's		C#	Python	
-			Ruby	

Recent languages: Scala, Clojure, F#, Lua, Go, ...

Others?

Programming language features

	Traditional, "static"	Static o-o	Scripting, "dynamic"	Mixed
Examples	C, Fortran	C++	Python, JS	Java, OCaml
Objects?	No	yer	yen	Yeu
Automatic	-1	۸/	110	Uo.
mem. mgt.?	No	No	yen	Ju
Static types?	Yee	yen	No	yee
Tagged values?	No	No	Yec	yes and vo

Down to details...

- MP0 "due" Thursday 9:30am
 - Install OCaml, write simple programs (approx. 1 hr. wo
 - Not graded, but uses handin
- MP1 due Tuesday, 1/22, 9:30am
 - Recursive functions on lists in OCaml
- First month of class is on lexing, parsing, type-checking, abstract syntax
 - Programming assignments in OCaml
 - Next two classes: intro to OCaml

OCaml

- OCaml is a popular functional programming languates
 which means
 - Calculations performed mainly by writing recursive functions on lists and trees
 - Automatic memory management
 - No assignable variables variables assigned once, never assigned; no while loops
 - Higher-order functions (but not until later in semester)
- Also, OCaml makes it very convenient to define a manipulate abstract syntax trees, which are of cru importance to us in this class.

Interactive use of OCaml

Like Python and Ruby, OCaml can be used by entering expressions function definitions interactively, in a "read-eval-print" loop:

```
~$ ocaml
# let pi = 3.14159;;
val pi : float = 3.14159
# let circum r = 2.0 *. pi *. r;; (* Use *. for floating pt mu
val circum : float -> float = <fun>
# circum 8.0;;
-: float = 50.26544
# let fac x = if x=0 then 1 else x * fac (x-1);;
Error: Unbound value fac
# let rec fac x = if x=0 then 1 else x * fac (x-1);
val fac : int -> int = <fun>
# fac 4;;
-: int = 24
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```

Interactive use of OCaml (cont.)

- Function application syntax: justaposition
- Variable and function names must begin with lower-c letter or underscore; can contain digits or apostrophe.
- No "return" statement: function body is an expression, a its value is returned.

```
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```

Interactive use of OCaml (cont.)

- if-then-else is a conditional expression, not a statemelike C's conditional expression (condition? expr : expr).
- #use must be used with care: if you load a file, then use definitions to define some other functions, and then relative the original file, the functions you defined in the meant will still use the old definitions.
- Although no types are given, OCaml is a statically ty language, like C or Java, unlike Python or JavaScript:

```
# circum 4;;
Error: This expression has type int but an expression was expected of ty
float
```

Types in OCaml

- OCaml provides powerful built-in types:
 - Primitive types: int, float, bool, char, string
 - Structured types:
 - Homogeneous lists: τ list, where τ is any type
 - Heterogeneous tuples: $\tau_1 * \tau_2 * \cdots * \tau_n$
 - Functions: $\tau_1 \rightarrow \tau_2 \rightarrow \cdots \rightarrow \tau_n$
- The class facility allows the definition of new types, as Java. However, we will not use classes. There is anot way to define new types — using the type facility — that will use heavily. We will discuss that next week.

Primitive types

- Primitive types int, float, bool, char, string are more less what you would expect:
 - int: 3, 74, -853
 - float: 3.0, .012
 - bool: true, false
 - char: 'a', '\n'
 - string: "sam I am"
- with one quirk: arithmetic operations for floats us period (+., *., etc.)

Primitive types (cont.)

- Comparison operations: =, <, >, <>, <=, >=
- Boolean operations: or or ||, & or &&, not
- String operations: ^ (concatenation); s.[i] (subscript)
- No automatic type conversions; use string_of_i int_of_string, float_of_int, etc.
- Various "modules" provide additional operations; th are loaded by entering open module-name;;.
 - String and Str modules have operations on strings.
 - Pervasives is a module that is always loaded don't need to open it.
 - Index of modules is linked at the bottom of the on manual — see "resources" tab in course web page.

Tuples

- Create "struct" by putting expressions in parenthe separated by commas:
 - (3, 5.0): int * float
 - (3, "abc", true): int * string * bool
- Exercise fill in types:
 - ('a', 'b'):
 - ('a', "a", 'a'):
 - (5, ("a", 'a')):
- Exercise create a value of the given type:
 - (3, (4, 5,0), "6") : int * (int * float) * string

Tuples (cont.)

- Use functions fst and snd to get elements of a pair. C works for pairs. (We'll see how to deal with bigger tuple next class.)
- Exercise: Write a function to add the elements of an int * pair:

```
• let addelts p = \int s1 p + 5nd p
```

```
addelts (3,4);; (* returns 7 *)
```

Lists

- Create a list by putting expressions, all of same type, square brackets separated by semicolons:
 - [3]: int list
 - [true; false; true]: bool list
 - [['a']; ['b'; 'c'; 'd']; []]: (char list) list
 - [3; 4.0]: *type error*
 - [3; [4]]: *type error*

Exercise — fill in types (or flag error):

```
• ['a'; 'b']: char list

• ['a'; 'b'; "c"]: type error

• [4; int_of_string "34"]: mt list

• [[4]; [5]; []]: (mt list) list

• [(1, 2); (3, 4)]: (mt * mt) list

• [(1, 2); (3, 4, 5)]: type error

• [(1, [3]); (4, [5; 6])]: (mt * (mt list)) list
```

Exercise — create a value of the given type (other than empty list):

• [3] : int list

• [3]] : (int list) list

• [(3 "3")] : (int * string) list

• ["3"] : (string list) list

• [(3,["3"])] : (int * string list) list

• [[(3, ["3"])]] : ((int * string list) list) list

Use @ to concatenate lists, :: to "cons" to start of list; least module to get functions hd, tl, nth, length, and other.

```
# open List;;
# let lis1 = [1; 2; 3];;
val lis1 : int list = [1; 2; 3]
# hd lis1;;
- : int = 1
# tl lis1;;
- : int list = [2; 3]
# 4 :: lis1;; (* N.B. non-destructive cons *)
- : int list = [4; 1; 2; 3]
# [4] @ lis1;; (* N.B. non-destructive append *)
- : int list = [4; 1; 2; 3]
# length lis1;;
- : int = 3
# nth lis1 2;;
- : int = 3
```

- Exercise: Write a function to compute the sum of the f two elements of an int list: addfirsttwo [5; 3; 2; 6] = You can assume the list is of length at least 2:
 - let addfirsttwo lis = hd lis + hd (tl lis)
- Exercise: Write a function to compute the sum the lengths of the first two elements of an (int l list: addfirsttwolengths [[5; 3]; [2]; [6; 2; 5; 3]]
 3. You can assume the list is of length at least 2:
 - let addfirsttwolengths lis = length (hd lis) + length (hd (tl lis))

Polymorphic functions

- OCaml detects the types of functions, so that variable a function types don't have to be declared.
- But consider:

```
let revpair p = (snd p, fst p)
```

- It is legal to write
 - revpair (3, 4.0)
 - revpair ("abc", true)
 - etc.
- In fact, for any types au and au',

revpair:
$$\tau * \tau' \rightarrow \tau' * \tau$$

Polymorphic functions (cont.)

OCaml realizes this and assigns revpair a polymorphic t
 a type with "type variables" in it:

```
# let revpair p = (snd p, fst p);;
val revpair : 'a * 'b -> 'b * 'a = <fun>
```

Read this as "revpair has type $\alpha * \beta \rightarrow \beta * \alpha$."

• Similarly, a function may operate on lists of any type. The length has type lpha list ightarrow int.

Polymorphic functions (cont.)

Ex: Write the polymorphic types of the following functions. (You can write elastic, 'b, etc. or α , β , as you prefer.)

let pair2list p = [fst p, snd p] $\alpha + \alpha \rightarrow \alpha \quad \text{let}$

addfirsttwolengths (defined above)

a list list - int

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Wrap-up

Today we discussed:

- What CS421 is about
- Basic OCaml programming and use of the top level

We discussed it because:

- It's good to know why you're learning this stuff
- We'll be using OCaml for MPs this semester. And one goal of the of is for you to learn functional programming.

In the next class, we will:

Talk more about OCaml — esp. writing functions on lists.

What to do now:

 MPO — not graded, but using OCaml a little will give you a much be feel for it than listening to me talk about it.