

# Exercise: simple expression evaluation

```
type exp = Operation of exp * binary_operation * exp
          | Id of string | Integer of int
and binary_operation = Equal | LessThan | Plus

type value = Int of int | Bool of bool

let rec eval e dict =
    (see midterm 1  
answer sheet)

and apply bop v1 v2 =
```

# eval for MJ

```
type value = IntV of int | StringV of string | BoolV of bool | NullV  
and state = (varname * value) list  
and varname = string  
  
let rec eval (e:exp) (sigma:state) (prog:program) : value = match e with  
  
  Null -> NullV  
  | True -> BoolV true  
  | False -> BoolV false  
  | Integer i -> IntV i  
  | String s -> StringV s  
  
  (* assume id is in state sigma *)  
  | Id id -> lookup id sigma
```

# applyOp for MJ (cont.)

```
type value = IntV of int | StringV of string | BoolV of bool | NullV  
| Operation(e1, bop, e2) -> (* for non-boolean operations *)  
    applyOp bop (eval e1 sigma prog) (eval e2 sigma prog)  
  
let applyOp (bop:binary_operation) (v1:value) (v2:value) : value =  
  match bop with  
  Multiplication -> match (v1,v2) with  
    (IntV i1, IntV i2) -> IntV (i1 * i2)  
  
  Plus -> match (v1,v2) with  
    (IntV i1, IntV i2) -> IntV (i1 + i2)  
    | (StringV s1, StringV s2) -> StringV (s1 ^ s2)  
    | : ;
```

# eval for MJ, with exceptions

```
type value = IntV of int | StringV of string | BoolV of bool | NullV  
and state = (varname * value) list  
and varname = string  
exception TypeError of string  
exception RuntimeError of string
```

| Id id -> if isIn id sigma  
then lookup id sigma  
else raise TypeError

| Not e -> match (eval e sigma prog) with  
BoolV b -> BoolV (not b)  
| \_ -> raise TypeError

# Ex: SOS for binary operations

(BINOPINT)  $e_1 + e_2, \sigma, \pi \Downarrow \text{IntV } (i_1 + i_2)$

$e_1, \sigma, \pi \Downarrow \text{IntV } i_1$

$e_2, \sigma, \pi \Downarrow \text{IntV } i_2$

(BINOPINT)  $e_1 * e_2, \sigma, \pi \Downarrow \text{IntV } (i_1 * i_2)$

$e_1, \sigma, \pi \Downarrow \text{IntV } i_1$

$e_2, \sigma, \pi \Downarrow \text{IntV } i_2$

(LESSTHAN)  $e_1 < e_2, \sigma, \pi \Downarrow \text{BoolV } (i_1 < i_2)$

$e_1, \sigma, \pi \Downarrow \text{IntV } i_1$

$e_2, \sigma, \pi \Downarrow \text{IntV } i_2$

# Ex: SOS for boolean operations

(ORTRUE)  $e_1 \mid\mid e_2, \sigma, \pi \Downarrow \text{BoolV true}$   
 $e_1, \sigma, \pi \Downarrow \text{BoolV true}$

(ORFALSE)  $e_1 \mid\mid e_2, \sigma, \pi \Downarrow \text{BoolV } t$   
 $e_1, \sigma, \pi \Downarrow \text{BoolV false}$   
 $e_2, \sigma, \pi \Downarrow \text{BoolV } t$

(ANDFALSE)  $e_1 \& \& e_2, \sigma, \pi \Downarrow \text{BoolV false}$   
 $e_1, \sigma, \pi \Downarrow \text{BoolV false}$

(ANDTRUE)  $e_1 \& \& e_2, \sigma, \pi \Downarrow \text{BoolV } t$   
 $e_1, \sigma, \pi \Downarrow \text{BoolV true}$   
 $e_2, \sigma, \pi \Downarrow \text{BoolV } t$

(NOT)  $\neg e, \sigma, \pi \Downarrow \text{BoolV } (\text{not } b)$   
 $e, \sigma, \pi \Downarrow \text{BoolV } \perp$

# Statements

- You will also need to write function `exec: statement → state`  
→ program → state to execute some simple statements:

```
statement = Block of (statement list)
| If of exp * statement * statement
| Assignment of id * exp
```

```
let rec exec s sigma prog = match s with
Assignment(s, e) ->
```

asgn = (eval e sigma prog) sigma

| If(e,s1,s2) -> match (eval e sigma prog) with  
BoolV b -> if b then exec s1 sigma prog  
else exec s2 sigma prog

| - → raise Type Error

## Method calls

eval clause for method call : *argument list*

| Method call (-, m, [e<sub>1</sub>; ...; e<sub>n</sub>])

① Look up m; suppose it is :

$$\left. \begin{array}{l} t \ m(x_1, \dots, x_n) \\ y_1, \dots, y_m; // \text{local vars} \\ s_1; \dots; s_n \\ \text{return } e_i \end{array} \right\}$$

② Evaluate eval e<sub>1</sub> sigma prog,  
eval e<sub>2</sub> sigma prog, ..., yielding  
values v<sub>1</sub>, ..., v<sub>n</sub>

③ Create state pairing x<sub>1</sub> with v<sub>1</sub>, x<sub>2</sub> with v<sub>2</sub>, etc.,  
and pairing y<sub>1</sub> with NullV, y<sub>2</sub> with NullV, etc.  
(call this state σ)

④ Execute statements s<sub>1</sub>, ..., s<sub>n</sub>, starting in  
state σ, yielding state σ'

⑤ Evaluate e<sub>i</sub> in σ' and return its value