Programming Languages and Compilers (CS 421)



Elsa L Gunter 2112 SC, UIUC

http://courses.engr.illinois.edu/cs421

Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha

12/1/15



Axiomatic Semantics

- Also called Floyd-Hoare Logic
- Based on formal logic (first order predicate calculus)
- Axiomatic Semantics is a logical system built from axioms and inference rules
- Mainly suited to simple imperative programming languages

12/1/15



Axiomatic Semantics

Used to formally prove a property (postcondition) of the state (the values of the program variables) after the execution of program, assuming another property (pre-condition) of the state holds before execution

12/1/15



Axiomatic Semantics

- Goal: Derive statements of form {P} C {Q}
 - P , Q logical statements about state,
 P precondition, Q postcondition,
 C program
- Example: $\{x = 1\} x := x + 1 \{x = 2\}$

12/1/15 4



Axiomatic Semantics

 Approach: For each type of language statement, give an axiom or inference rule stating how to derive assertions of form {P} C {Q}

where C is a statement of that type

 Compose axioms and inference rules to build proofs for complex programs

12/1/15



5

Axiomatic Semantics

- An expression {P} C {Q} is a partial correctness statement
- For total correctness must also prove that C terminates (i.e. doesn't run forever)
 - Written: [P] C [Q]
- Will only consider partial correctness here



Language

We will give rules for simple imperative language

<command>

- ::= <variable> := <term>
 - <command>; ...; <command>
 - if <statement> then <command> else
- <command>
- | while <statement> do <command>
- Could add more features, like for-loops

12/1/15



Substitution

- Notation: P[e/v] (sometimes P[v <- e])</p>
- Meaning: Replace every v in P by e
- Example:

$$(x + 2) [y-1/x] = ((y - 1) + 2)$$

12/1/15



The Assignment Rule

$${P [e/x]} x := e {P}$$

Example:
$$\frac{}{\{ ? \} x := y \{x = 2\}}$$

12/1/15

The Assignment Rule

$${P [e/x]} x := e {P}$$

Example:

$$\{ = 2 \} x := y \{ x = 2 \}$$

12/1/15 10



The Assignment Rule

$${P [e/x]} x := e {P}$$

Example:

$${y=2} x := y {x=2}$$

12/1/15

The Assignment Rule

$${P [e/x]} x := e {P}$$

Examples:

$${y = 2} x := y {x = 2}$$

$$\overline{\{y = 2\} \ x := 2 \ \{y = x\}}$$

$${x + 1 = n + 1} x := x + 1 {x = n + 1}$$

$${2 = 2} x := 2 {x = 2}$$

12/1/15

11



The Assignment Rule - Your Turn

• What is the weakest precondition of

$$x := x + y \{x + y = w - x\}?$$

{ ? ? }
$$x := x + y$$
 $\{x + y = w - x\}$

12/1/15

The A

The Assignment Rule – Your Turn

What is the weakest precondition of

$$x := x + y \{x + y = w - x\}$$
?

$$\{(x + y) + y = w - (x + y)\}$$

 $x := x + y$
 $\{x + y = w - x\}$

12/1/15 14



Precondition Strengthening

$$\frac{P \Rightarrow P' \quad \{P'\} C \{Q\}}{\{P\} C \{Q\}}$$

- Meaning: If we can show that P implies P' (P→ P') and we can show that {P'} C {Q}, then we know that {P} C {Q}
- P is stronger than P' means P → P'

12/1/15



13

15

Precondition Strengthening

Examples:

$$\frac{x = 3 \Rightarrow x < 7 \ \{x < 7\} \ x := x + 3 \ \{x < 10\}}{\{x = 3\} \ x := x + 3 \ \{x < 10\}}$$

True
$$\Rightarrow$$
 2 = 2 {2 = 2} x:= 2 {x = 2}
{True} x:= 2 {x = 2}

$$\frac{x=n \Rightarrow x+1=n+1 \quad \{x+1=n+1\} \; x:=x+1 \; \{x=n+1\}}{\{x=n\} \; x:=x+1 \; \{x=n+1\}}$$

12/1/15 16



Which Inferences Are Correct?

$$\frac{\{x > 0 \& x < 5\} \ x := x * x \{x < 25\}}{\{x = 3\} \ x := x * x \{x < 25\}}$$

$$\frac{\{x * x < 25\} x := x * x \{x < 25\}}{\{x > 0 \& x < 5\} x := x * x \{x < 25\}}$$

12/1/15

4

Which Inferences Are Correct?

$$\frac{\{x > 0 \& x < 5\} \ x := x * x \{x < 25\}}{\{x = 3\} \ x := x * x \{x < 25\}}$$

$$\frac{\{x = 3\} \times := x * x \{x < 25\}}{\{x > 0 \& x < 5\} \times := x * x \{x < 25\}}$$

$$\frac{\{x * x < 25\} x := x * x \{x < 25\}}{\{x > 0 \& x < 5\} x := x * x \{x < 25\}}$$

12/1/15



Sequencing

$${P} C_1 {Q} {Q} C_2 {R}$$

 ${P} C_1; C_2 {R}$

Example:

$${z = z \& z = z} x := z {x = z \& z = z}$$

 ${x = z \& z = z} y := z {x = z \& y = z}$
 ${z = z \& z = z} x := z; y := z {x = z & y = z}$

12/1/15



$${P} C_1{Q} {Q} C_2{R}$$

 ${P} C_1; C_2{R}$

Example:

$${z = z & z = z} x := z {x = z & z = z}$$

 ${x = z & z = z} y := z {x = z & y = z}$
 ${z = z & z = z} x := z; y := z {x = z & y = z}$

12/1/15 20



Postcondition Weakening

$$\begin{array}{cccc} \{P\} C \{Q'\} & Q' & \longrightarrow Q \\ \hline \{P\} C \{Q\} & \end{array}$$

Example:

12/1/15



Rule of Consequence

$$P \xrightarrow{P'} P' \quad \{P'\} C \{Q'\} \quad Q' \xrightarrow{P} Q$$
$$\{P\} C \{Q\}$$

- Logically equivalent to the combination of Precondition Strengthening and Postcondition Weakening
- Uses P → P and Q → Q

12/1/15 22



If Then Else

Example: Want

Suffices to show:

- (1) ${y=a&x<0}$ $y:=y-x {y=a+|x|}$ and
- (4) $\{y=a¬(x<0)\}\ y:=y+x\ \{y=a+|x|\}$

12/1/15

23

19

21



${y=a&x<0} y:=y-x {y=a+|x|}$

- (3) (y=a&x<0)**→** y-x=a+|x|
- (2) $\{y-x=a+|x|\}\ y:=y-x\ \{y=a+|x|\}$
- (1) y=a&x<0 y:=y-x y=a+|x|
- (1) Reduces to (2) and (3) by Precondition Strengthening
- (2) Follows from assignment axiom
- (3) Because $x<0 \rightarrow |x| = -x$



${y=a¬(x<0)} y:=y+x {y=a+|x|}$

- (6) $(y=a¬(x<0)) \rightarrow (y+x=a+|x|)$
- (5) $\{y+x=a+|x|\}\ y:=y+x\ \{y=a+|x\}\}$
- (4) $\{y=a¬(x<0)\}\ y:=y+x\ \{y=a+|x|\}$
- (4) Reduces to (5) and (6) by Precondition Strengthening
- (5) Follows from assignment axiom
- (6) Because $not(x<0) \rightarrow |x| = x$

12/1/15

25



- (1) ${y=a&x<0}y:=y-x{y=a+|x|}$
- (4) ${y=a¬(x<0)}y:=y+x{y=a+|x|}$

if x < 0 then y:= y-x else y:= y+x $\{y=a+|x|\}$

By the if then else rule

12/1/15 26



While

- We need a rule to be able to make assertions about while loops.
 - Inference rule because we can only draw conclusions if we know something about the body
 - Let's start with:

12/1/15

27



While

The loop may never be executed, so if we want P to hold after, it had better hold before, so let's try:

12/1/15 28



While

- If all we know is P when we enter the while loop, then we all we know when we enter the body is (P and B)
- If we need to know P when we finish the while loop, we had better know it when we finish the loop body:

$$\frac{ \{ P \text{ and B} \} \ C \ \{ P \}}{ \{ P \} \ \text{while} \ B \ \text{do} \ C \ \{ P \}}$$

12/1/15

29



While

- We can strengthen the previous rule because we also know that when the loop is finished, not B also holds
- Final while rule:

12/1/15



While

 P satisfying this rule is called a loop invariant because it must hold before and after the each iteration of the loop

12/1/15 31



While

- While rule generally needs to be used together with precondition strengthening and postcondition weakening
- There is NO algorithm for computing the correct P; it requires intuition and an understanding of why the program works

12/1/15 32



Example

Let us prove

$$\{x>= 0 \text{ and } x = a\}$$

fact := 1;
while $x > 0$ do (fact := fact * x; x := x -1)
 $\{fact = a!\}$

12/1/15



Example

 We need to find a condition P that is true both before and after the loop is executed, and such that

(P and not
$$x > 0$$
) \rightarrow (fact = a!)

12/1/15 34



Example

First attempt:

$${a! = fact * (x!)}$$

- Motivation:
- What we want to compute: a!
- What we have computed: fact
 which is the sequential product of a down through (x + 1)
- What we still need to compute: x!

12/1/15 35



33

Example

By post-condition weakening suffices to show



Problem

- {a! = fact * (x!) and not (x > 0)} → {fact = a!}
- Don't know this if x < 0
- Need to know that x = 0 when loop terminates
- Need a new loop invariant
- Try adding $x \ge 0$
- Then will have x = 0 when loop is done

12/1/15

37

39

41



Second try, combine the two:

$$P = \{a! = fact * (x!) and x >= 0\}$$

Again, suffices to show

1.
$$\{x > = 0 \text{ and } x = a\}$$

fact := 1;

while
$$x > 0$$
 do (fact := fact * x; $x := x - 1$)

 $\{P \text{ and not } x > 0\}$

and

2. $\{P \text{ and not } x > 0\} \rightarrow \{fact = a!\}$

12/1/15 38



Example

For 2, we need

$${a! = fact * (x!) and x >= 0 and not (x > 0)}$$

But
$$\{x >= 0 \text{ and not } (x > 0)\} \rightarrow \{x = 0\} \text{ so}$$

$$fact * (x!) = fact * (0!) = fact$$

Therefore

{a! = fact * (x!) and x >= 0 and not (x > 0)}
$$\rightarrow$$
 {fact = a!}

12/1/15



Example

For 1, by the sequencing rule it suffices to show

3.
$$\{x \ge 0 \text{ and } x = a\}$$

$${a! = fact * (x!) and x >= 0}$$

And

4.
$$\{a! = fact * (x!) and x >= 0\}$$

while
$$x > 0$$
 do

$$(fact := fact * x; x := x - 1)$$

$$\{a! = \text{fact } * (x!) \text{ and } x >= 0 \text{ and not } (x > 0)\}$$

12/1/15 40



Example

Suffices to show that

$${a! = fact * (x!) and x >= 0}$$

holds before the while loop is entered and that if

$$\{(a! = fact * (x!)) \text{ and } x >= 0 \text{ and } x > 0\}$$

holds before we execute the body of the loop, then

$$\{(a! = fact * (x!)) \text{ and } x >= 0\}$$

holds after we execute the body

12/1/15



Example

By the assignment rule, we have

$${a! = 1 * (x!) and x >= 0}$$

$${a! = fact * (x!) and x >= 0}$$

Therefore, to show (3), by precondition strengthening, it suffices to show

$$(x>= 0 \text{ and } x = a) \rightarrow$$

$$(a! = 1 * (x!) and x >= 0)$$

12/1/15



Example

$$(x>= 0 \text{ and } x = a) \rightarrow$$

 $(a! = 1 * (x!) \text{ and } x >= 0)$
holds because $x = a \rightarrow x! = a!$

Have that $\{a! = fact * (x!) and x >= 0\}$ holds at the start of the while loop

12/1/15

Example

To show (4): $\{a! = \text{fact} * (x!) \text{ and } x >= 0\}$ while x > 0 do $\{\text{fact} := \text{fact} * x; x := x - 1\}$ $\{a! = \text{fact} * (x!) \text{ and } x >= 0 \text{ and not } (x > 0)\}$ we need to show that $\{(a! = \text{fact} * (x!)) \text{ and } x >= 0\}$ is a loop invariant

12/1/15 44



Example

We need to show:

$$\{(a! = fact * (x!)) \text{ and } x \ge 0 \text{ and } x \ge 0\}$$

(fact = fact * x; x := x - 1)
 $\{(a! = fact * (x!)) \text{ and } x \ge 0\}$

We will use assignment rule, sequencing rule and precondition strengthening

12/1/15



43

45

47

Example

By the assignment rule, we have
$$\{(a! = \text{fact} * ((x-1)!)) \text{ and } x - 1 \ge 0\}$$

$$x := x - 1$$

$$\{(a! = \text{fact} * (x!)) \text{ and } x \ge 0\}$$
 By the sequencing rule, it suffices to show
$$\{(a! = \text{fact} * (x!)) \text{ and } x \ge 0 \text{ and } x \ge 0\}$$

$$\text{fact} = \text{fact} * x$$

$$\{(a! = \text{fact} * ((x-1)!)) \text{ and } x - 1 \ge 0\}$$

12/1/15 46



Example

12/1/15



Example

However

fact * x * (x - 1)! = fact * x
and (x > 0)
$$\rightarrow$$
 x - 1 >= 0
since x is an integer,so
{(a! = fact * (x!)) and x >= 0 and x > 0} \rightarrow
{(a! = (fact * x) * ((x-1)!)) and x - 1 >= 0}



Example

```
Therefore, by precondition strengthening \{(a! = fact * (x!)) \text{ and } x >= 0 \text{ and } x > 0\}
fact = fact * x
\{(a! = fact * ((x-1)!)) \text{ and } x - 1 >= 0\}
```

This finishes the proof

12/1/15