

Hoare Logic 2

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Hoare Triple

$P \{ \dots \text{code} \dots \} Q$

$$P[e/x] \{ x := e \} P$$
$$P \{ C_1 \} R \quad R \{ C_2 \} Q$$

$$P \{ C_1; C_2 \} Q$$
$$P \wedge b \{ C_1 \} Q \quad P \wedge \neg b \{ C_2 \} Q$$

$$P \{ \text{if } b \text{ then } C_1 \text{ else } C_2 \} Q$$

While Rule

$$\frac{P \wedge b \{ C \} P}{P \{ \text{While } b \ C \} P \wedge \neg b}$$

(**P is a loop invariant**)

Rule of Consequence

$$\frac{P \rightarrow P' \quad P' \{ C \} Q' \quad Q' \rightarrow Q}{P \{ C \} Q}$$

Sample Proofs

- sum of n
- fibonacci
- list append
- list reverse
- termination

Sum of n

$x = 0 \ \& \ y = 0$

{

While $y < n$

$y := y + 1;$

$x := x + y$

}

$x = 1 + \dots + n$

$P \equiv x = 1 + \dots + y \wedge y \leq n$

| | |
|--|---|
| $x = 0 \wedge y = 0 \rightarrow x = 1 + \dots + y \wedge y \leq n$ | ✓ |
| $x = 1 + \dots + y \wedge y \leq n \wedge \neg(y < n) \rightarrow x = 1 + \dots + n$ | ✓ |
| $x = 1 + \dots + y \wedge y \leq n \wedge y < n \rightarrow$? | ✓ |

$x + y + 1 = 1 + \dots + y + 1 \wedge y + 1 \leq n$

$$\begin{array}{l} \{y := y + 1\} \quad x + y = 1 + \dots + y \wedge y \leq n \\ \{x := x + y\} \quad x = 1 + \dots + y \wedge y \leq n \end{array}$$

?

$$\{y := y + 1; x := x + y\} \quad x = 1 + \dots + y \wedge y \leq n$$

$$x = 1 + \dots + y \wedge y \leq n \wedge y < n \quad \{y := y + 1; x := x + y\} \quad x = 1 + \dots + y \wedge y \leq n$$

$$x = 1 + \dots + y \wedge y \leq n \quad \{\text{While } y < n \dots\} \quad x = 1 + \dots + y \wedge y \leq n \wedge \neg(y < n)$$

$$x = 0 \wedge y = 0 \quad \{\text{While } \dots\} \quad x = 1 + \dots + n$$

Fibonacci

$x = 0 \ \& \ y = 1 \ \& \ z = 1 \ \& \ 1 \leq n$

{

While $z < n$

$y := x + y;$

$x := y - x;$

$z := z + 1$

}

$y = \text{fib } n$

$P \equiv y = \text{fib } z \ \wedge \ x = \text{fib } (z-1)$

$\wedge \ z \leq n$

$$\begin{aligned}
 & x = 0 \wedge y = 1 \wedge z = 0 \wedge 1 \leq n \rightarrow y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \quad \checkmark \\
 & y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge \neg(z < n) \rightarrow y = \text{fib } n \quad \checkmark \\
 & y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge z < n \rightarrow \quad ? \quad \checkmark
 \end{aligned}$$

$$\boxed{x+y = \text{fib } (z+1) \wedge x+y-x = \text{fib } (z+1-1) \wedge z + 1 \leq n}$$

$$\begin{aligned}
 \{y := x + y\} & \quad y = \text{fib } (z+1) \wedge y-x = \text{fib } (z+1-1) \wedge z + 1 \leq n \\
 \{x := y - x\} & \quad y = \text{fib } (z+1) \wedge x = \text{fib } (z+1-1) \wedge z + 1 \leq n \\
 \{z := z + 1\} & \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n
 \end{aligned}$$

$$? \quad \{y := x + y; x := y - x; z := z + 1\} \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n$$

$$y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge z < n \quad \{y := x + y; x := y - x; z := z + 1\} \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n$$

$$y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \quad \{\text{While } z < n \dots\} \quad y = \text{fib } z \wedge x = \text{fib } (z-1) \wedge z \leq n \wedge \neg(z < n)$$

$$x = 0 \wedge y = 1 \wedge z = 0 \wedge 1 \leq n \quad \{\text{While } \dots\} \quad y = \text{fib } n$$

List length

$x = \text{lst} \ \& \ y = 0$

$P \equiv \text{len lst} = y + \text{len } x$

{

 While $x \neq []$

$x := \text{tl } x;$

$y := y + 1$

}

$y = \text{len lst}$

| | |
|--|---|
| $x = \text{lst} \wedge y = 0 \rightarrow \text{len lst} = y + \text{len } x$ | ✓ |
| $\text{len lst} = y + \text{len } x \wedge \neg(x \neq []) \rightarrow y = \text{len lst}$ | ✓ |
| $\text{len lst} = y + \text{len } x \wedge x \neq [] \rightarrow ?$ | ✓ |

$$\boxed{\text{len lst} = y + 1 + \text{len}(\text{tl } x)}$$

$$\begin{array}{ll} \{x := \text{tl } x\} & \text{len lst} = y + 1 + \text{len } x \\ \{y := y + 1\} & \text{len lst} = y + \text{len } x \end{array}$$

?

$$\{x := \text{tl } x; y := y + 1\} \quad \text{len lst} = y + \text{len } x$$

$$\text{len lst} = y + \text{len } x \wedge x \neq [] \quad \{x := \text{tl } x; y := y + 1\} \quad \text{len lst} = y + \text{len } x$$

$$\text{len lst} = y + \text{len } x \quad \{\text{While } x \neq [] \dots\} \quad \text{len lst} = y + \text{len } x \wedge \neg(x \neq [])$$

$$x = \text{lst} \wedge y = 0 \quad \{\text{While } \dots\} \quad y = \text{len lst}$$

List reverse

$x = \text{lst} \ \& \ y = []$ $P \equiv \text{lst} = \text{rev } y \ @ \ x$

{

 While $x \neq []$

$y := \text{hd } x :: y;$

$x := \text{tl } x$

}

$y = \text{rev lst}$

| | |
|--|---|
| $x = \text{lst} \wedge y = [] \rightarrow \text{lst} = \text{rev } y @ x$ | ✓ |
| $\text{lst} = \text{rev } y @ x \wedge \neg(x \neq []) \rightarrow y = \text{rev } \text{lst}$ | ✓ |
| $\text{lst} = \text{rev } y @ x \wedge x \neq [] \rightarrow$? | ✓ |

$$\boxed{\text{lst} = \text{rev } (\text{hd } x @ y) @ (\text{tl } x)}$$

$$\begin{array}{ll} \{y := \text{hd } x @ y\} & \text{lst} = \text{rev } y @ (\text{tl } x) \\ \{x := \text{tl } x\} & \text{lst} = \text{rev } y @ x \end{array}$$

$$? \quad \{y := \text{hd } x @ y; x := \text{tl } x\} \quad \text{lst} = \text{rev } y @ x$$

$$\text{lst} = \text{rev } y @ x \wedge x \neq [] \quad \{y := \text{hd } x @ y; x := \text{tl } x\} \quad \text{lst} = \text{rev } y @ x$$

$$\text{lst} = \text{rev } y @ x \quad \{\text{While } x \neq [] \dots\} \quad \text{lst} = \text{rev } y @ x \wedge \neg(x \neq [])$$

$$x = \text{lst} \wedge y = [] \quad \{\text{While } \dots\} \quad y = \text{rev } \text{lst}$$