| Programming Languages and Compilers (CS 421) <br> Elsa L Gunter <br> 2112 SC, UIUC <br> http://www.cs.uiuc.edu/class/cs421/ <br> Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha <br> $8 / 31 / 11$ |  |
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## Recursion Example

```
Compute \(\mathrm{n}^{2}\) recursively using:
    \(\mathrm{n}^{2}=(2 * n-1)+(n-1)^{2}\)
\# let rec nthsq \(\mathrm{n}=\) (* rec for recursion *)
    match \(\mathrm{n} \quad\) (* pattern matching for cases *)
    with 0 -> 0 (* base case *)
    \(\mid \mathrm{n}->(2 * \mathrm{n}-1) \quad(*\) recursive case *)
        + nthsq ( \(n-1\) );i, (* recursive call *)
```

    val nthsq : int -> int = <fun>
    \# nthsq 3;
    : int =9
    Structure of recursion similar to inductive proof
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## Recursion and Induction

\# let rec nthsq $\mathrm{n}=$ match n with $0->0$

$$
\mid n->(2 * n-1)+n t h s q(n-1) ;
$$

- Base case is the last case; it stops the computation
- Recursive call must be to arguments that are somehow smaller - must progress to base case
- if or match must contain base case
- Failure of these may cause failure of termination


## Structural Recursion : List Example

\# let rec length list = match list
with [ ] -> 0 (* Nil case *)
| x :: xs -> 1 + length xs;;; (* Cons case *)
val length : 'a list -> int = <fun>
\# length [5; 4; 3; 2];,

- : int = 4
- Nil case [ ] is base case
- Cons case recurses on component list xs

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## Forward Recursion

- In structural recursion, you split your input into components
- Forward recursion limited form of structural recursion
- In forward recursion:
- First call function recursively on all recursive components
- Recursive calls use only components and input agruments, no other computation
- Build result from results of component

| Forward Recursion: Examples |
| :--- |
| \# let rec double_up list = |
| match list |
| with [ ] -> [ ] |
| \| (x :: xs) -> (x :: x : : double_up xs);; |
| val double_up : 'a list -> 'a list = <fun> |
| \# let rec poor_rev list = |
| match list |
| with [] -> [] |
| \| (x::xs) -> poor_rev xs @ [x];i; |
| val poor_rev : 'a list -> 'a list = <fun> |
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## Mapping Recursion

- One common form of structural recursion applies a function to each element in the structure
\# let rec doubleList list = match list with [ ] -> [ ]
| x::xs -> 2 * x :: doubleList xs;,;
val doubleList : int list -> int list $=<$ fun $>$
\# doubleList [2;3;4];;
- : int list = $4 ; 6 ; 8]$


## Recall Map

\# let rec map f list =
match list
with [] -> []
| (h::t) -> (f h) :: (map ft);;
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>
\# map plus_two fib5;;

- : int list = [10; 7; 5; 4; 3; 3]
- Same as List.map


## Mapping Recursion

- Can use the higher-order recursive map function instead of direct recursion
\# let doubleList list =
List.map (fun x -> 2 * x) list;;
val doubleList : int list $->$ int list $=<$ fun>
\# doubleList $[2 ; 3 ; 4] ;$;
- : int list = 4 ; 6; 8]
- Same function, but no rec


## Folding Recursion

Another common form "folds" an operation over the elements of the structure
\# let rec multList list = match list with [ ] -> 1
| x::xs -> x * multList xs;;
val multList : int list $->$ int $=<$ fun $>$
\# multList [2;4;6];;

- : int = 48
- Computes (2 * (4*(6*1)))


## Folding Recursion

- multList folds to the right


## - Same as:

\# let multList list =
List.fold_right
(fun $\mathrm{p}->$ fun $\mathrm{x}->\mathrm{x}$ * p )
list $1 ;$
val multList : int list $->$ int $=$ <fun >
\# multList [2;4;6];;

- : int $=48$

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## Folding Functions over Lists

How are the following functions similar?
\# let rec sumlist list = match list with
[ ] -> 0 | $\mathrm{x}:$ :xs -> $\mathrm{x}+$ sumlist xs ;;
val sumlist : int list $->$ int $=<$ fun $>$
\# sumlist [2;3;4];;

- : int = 9
\# let rec prodlist list = match list with
[] -> 1 | x::xs -> x * prodlist xs;;
val prodlist : int list $->$ int $=<$ fun $>$
\# prodlist [2;3;4];i
- : int = 24


## Folding - Forward Recursion

\# let sumlist list = fold_right (+) list 0; ;
val sumlist : int list $->$ int $=$ <fun>
\# sumlist [2;3;4];,

- : int = 9
\# let prodlist list = fold_right (*) list 1;,
val prodlist : int list -> int $=$ <fun>
\# prodlist [2;3;4];;
- : int = 24


## Map from Fold

\# let map flist =
fold_right (fun x y -> fx:: y) list [ ];;
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>
\# map ((+)1) [1;2;3];;

- : int list = $2 ; 3 ; 4]$
- Can you write fold_right (or fold_left) with just map? How, or why not?


## Folding

\# let rec fold_left falist = match list with [] -> a | (x :: xs) -> fold_left f (f a x) xs;; val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a = <fun>
fold_left fa $\left[x_{1} ; x_{2} ; \ldots ; x_{n}\right]=f\left(\ldots\left(f\left(f_{\text {a }} x_{1}\right) x_{2}\right) . ..\right) x_{n}$
\# let rec fold_right $f$ list $b=$ match list
with [ ] -> b | (x :: xs) -> f x (fold_right f xs b);; val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b = <fun>
fold_right $\mathrm{f}\left[\mathrm{x}_{1} ; \mathrm{x}_{2} ; \ldots ; \mathrm{x}_{\mathrm{n}}\right] \mathrm{b}=\mathrm{f} \mathrm{x}_{1}\left(\mathrm{f} \mathrm{x}_{2}\left(\ldots\left(\mathrm{f} \mathrm{x}_{\mathrm{n}} \mathrm{b}\right) . ..\right)\right)$

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## Encoding Recursion with Fold

\# let rec append list1 list2 = match list1 with
[ ] -> list2 | x::xs -> x :: append xs list2;;
val append : 'a list -> 'a list -> a list = <fun>
Base Case Operation Recursive Call
\# let append list1 list2 =
fold_right (fun $x$ y -> x :: y) list1 list2;;
val append : 'a list -> 'a list -> 'a list = <fun>
\# append $[1 ; 2 ; 3][4 ; 5 ; 6]$;

- : int list = $1 ; 2 ; 3 ; 4 ; 5 ; 6]$

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## An Important Optimization

- When a function call is made,
 saved to the stack so we know to where to return when the call is finished
- What if $f$ calls $g$ and $g$ calls h, but calling $h$ is the last thing $g$ does (a tail cal/)?


## An Important Optimization

- When a function call is made, Tail the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if $f$ calls $g$ and $g$ calls $h$, but calling $h$ is the last thing $g$ does (a tail cal)?
- Then $h$ can return directly to $f$ instead of $g$


## Tail Recursion - Example

\# let rec rev_aux list revlist = match list with []-> revlist
| x :: xs -> rev_aux xs (x::revlist);;
val rev_aux : 'a list -> 'a list -> 'a list = <fun>
\# let rev list = rev_aux list [ ] ;
val rev : 'a list -> 'a list = <fun>

- What is its running time?


## Comparison

- $\operatorname{rev}[1,2,3]=$
- rev_aux [1,2,3] [ ] =
- rev_aux $[2,3][1]=$
- rev_aux [3] [2,1] =
- rev_aux [ ] [3,2,1] = [3,2,1]


## Tail Recursion

- A recursive program is tail recursive if all recursive calls are tail calls
- Tail recursive programs may be optimized to be implemented as loops, thus removing the function call overhead for the recursive calls
- Tail recursion generally requires extra "accumulator" arguments to pass partial results
- May require an auxiliary function

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## Comparison

- poor_rev $[1,2,3]=$
- (poor_rev [2,3]) @ [1] =
- ((poor_rev [3]) @ [2]) @ [1] =
- (((poor_rev [ ]) @ [3]) @ [2]) @ [1] =
- (([ ] @ [3]) @ [2]) @ [1]) =
- ([3] @ [2]) @ [1] =
- (3:: ([ ] @ [2])) @ [1] =
- $[3,2]$ @ $[1]=$
- $3::([2] @[1])=$
- 3 :: (2:: ([ ] @ [1])) = [3, 2, 1]


## Folding

- Can replace recursion by fold_right in any forward primitive recursive definition
- Primitive recursive means it only recurses on immediate subcomponents of recursive data structure
- Can replace recursion by fold_left in any tail primitive recursive definition

