## Problem Set 2

## Fall 11

Due: 27th September, 2011, 11:00 am before class begins
Please follow the homework format guidelines posted on the class web page:

> http://www.cs.illinois.edu/class/fa11/cs373/

Also, note that Problem 6 is an extra credit question.

1. [Category: NFA Comprehension, Points: 20]

Consider the following NFA $M$.

(a) Formally show that $M$ accepts the string $w_{1}=a b a a a b$ and string $w_{2}=b a b a a b$.
(b) Give a formal definition of the language that $M$ recognizes. Briefly describe why $M$ recognizes it.
2. [Category: NFA Construction, Points: 20]

Construct a non-deterministic finite automata that accepts the language $\{01,012\}^{*}$ over the alphabet, $\{0,1,2\}$. Your automata should contain only three states.

Hint: Think nondeterminism, and $\epsilon$ is your friend.
3. [Category: Construction, Points: 20]

For a string $w$, the reverse of $w$ is defined as the string obtained by reading $s$ from right to left, denoted by $w^{-1}$. For example, if $w=a b c$, then $w^{-1}=c b a$; if $w=a b a b$, then $w^{-1}=b a b a$.

For a language $L$, the reverse of $L$ is defined as the language

$$
\operatorname{reverse}(L)=\left\{w^{-1} \mid w \in L\right\}
$$

Let $A=\left\{Q, \Sigma, \delta, q_{0}, F\right\}$ be a DFA accepting $L$, construct an NFA $B$ with no more than $|Q|+1$ states that will accept reverse $(L)$. Give the formal definition of $B$ (in tuple notation, no diagram). You should also argue how/why this NFA works (intuitive explanation is enough).
4. [Category: Regular Expressions, Points: $4+4+4+8$ ]

Give a regular expression for each of the following languages; the alphabet is $\{a, b\}$.

- The set of all words that end with a $b$.
- The set of all words that begin with $a a$ and end with $a b$.
- The set of all words such that every occurrence of $a$ is immediately followed by a $b$.
- The set of all words such that the number of changes from $a$ to $b$ is the same as the number of changes from $b$ to $a$ when read left to right.
(E.g., aabbbabbbba is in the language, as there are two places where $a$ 's change to $b$ 's and two places where $b$ 's change to $a$ 's; however, $a a b b b a b$ is not in the language as $a$ 's change to $b$ 's twice, while $b$ 's change to $a$ 's only once).

5. [Category: NFA to DFA Conversion, Points: 20]

Convert the following NFA to a DFA using the subset construction, and show the state diagram.


You can check your answer (if you wish) by feeding a DFA to the website:

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http://pub.ist.ac.at/automata_tutor/solve?pid=16
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However, the site does not check if you are describing a DFA; also note that you will lose points if you do not follow the subset construction.
6. [Category: Extra Credit:, Points: 20]

Give a language $L$ over the alphabet $\Sigma=\{a, b\}$ such that any DFA accepting $L$ requires at least 3 final states. Prove that the language $L$ you give has this property.

