System Calls and I/O

Announcements

- cs241help-su12@cs.illinois.edu
  - Was misconfigured on Monday/Tuesday
  - If you sent an e-mail to that address, please resend it.
  - Tested, verified, and it's working now!
- Nightly Autograder
- HW1 Due Tonight (11:59pm, on svn)

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Three types of calls…

- Function Call
- Library Function Call
  - “Library Function”
  - “Library Call”
- System Call
  - “syscall”

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Three types of calls…

Function Calls

Library Function Calls

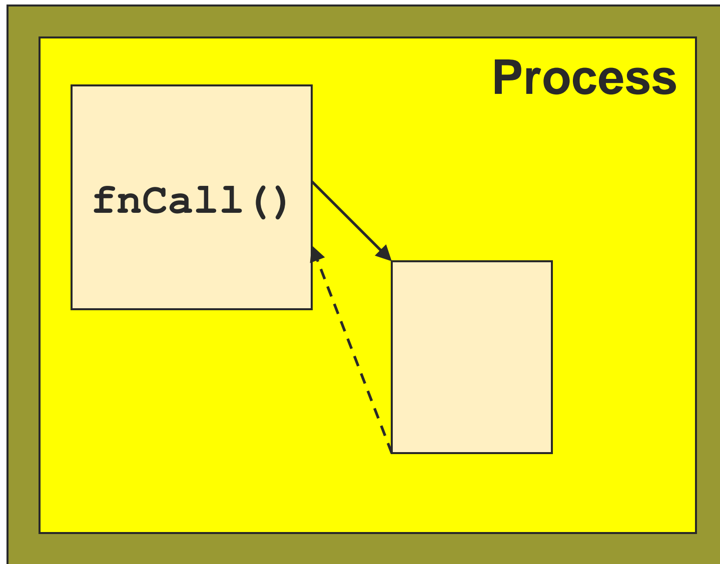
System Calls

Three types of calls…

- **Function Call**
- Library Function Call
  - "Library Function"
  - "Library Call"
- **System Call**
  - "syscall"

System Calls versus Function Calls

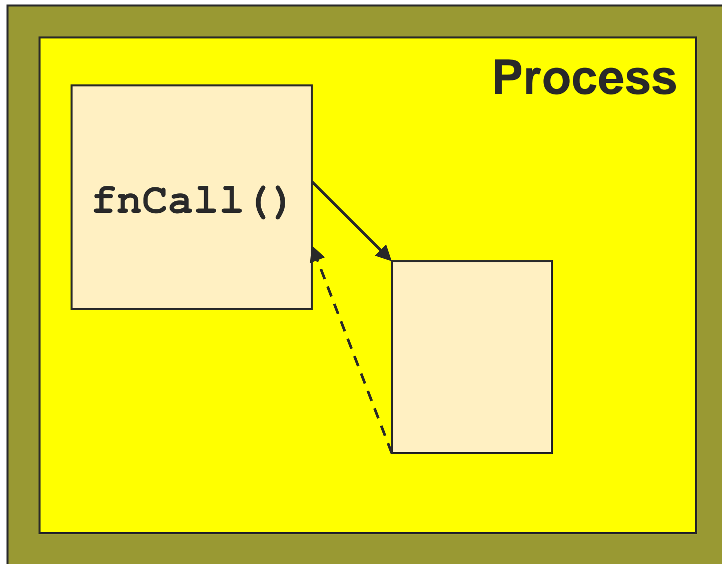
Function Call



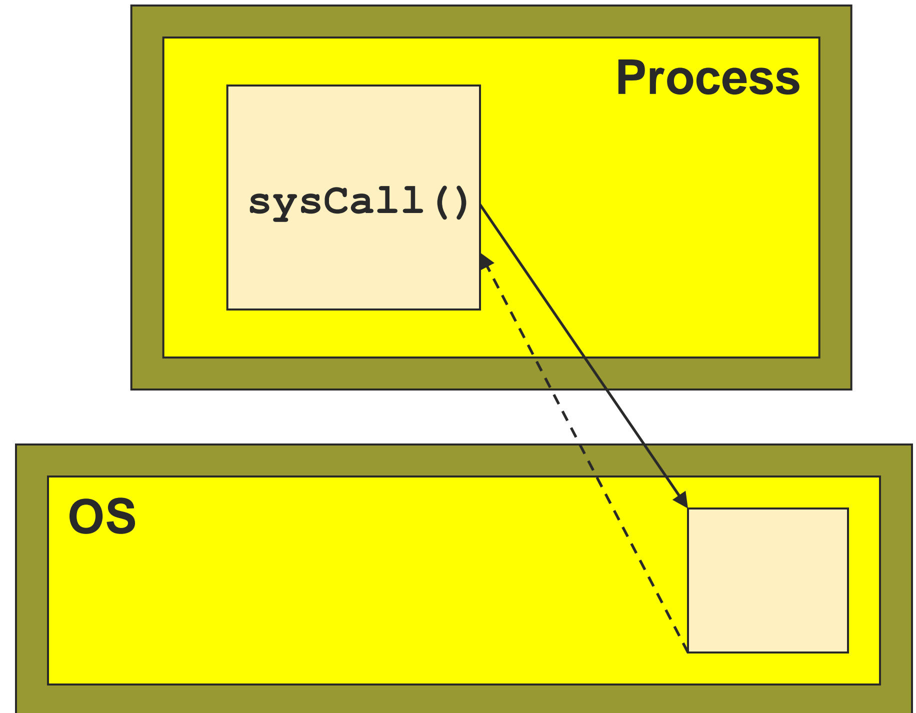
Caller and callee are in the same Process
- Same user
- Same "domain of trust"

System Calls versus Function Calls

Function Call



System Call



Caller and callee are in the same Process
- Same user
- Same "domain of trust"

- OS is trusted; user is not.
- OS has super-privileges; user does not
- Must take measures to prevent abuse

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System Calls

- System Calls
  - A request to the operating system to perform some activity
- System calls are expensive
  - The system needs to perform many things before executing a system call
    - The computer (hardware) saves its state
    - The OS code takes control of the CPU, privileges are updated.
    - The OS examines the call parameters
    - The OS performs the requested function
    - The OS saves its state (and call results)
    - The OS returns control of the CPU to the caller

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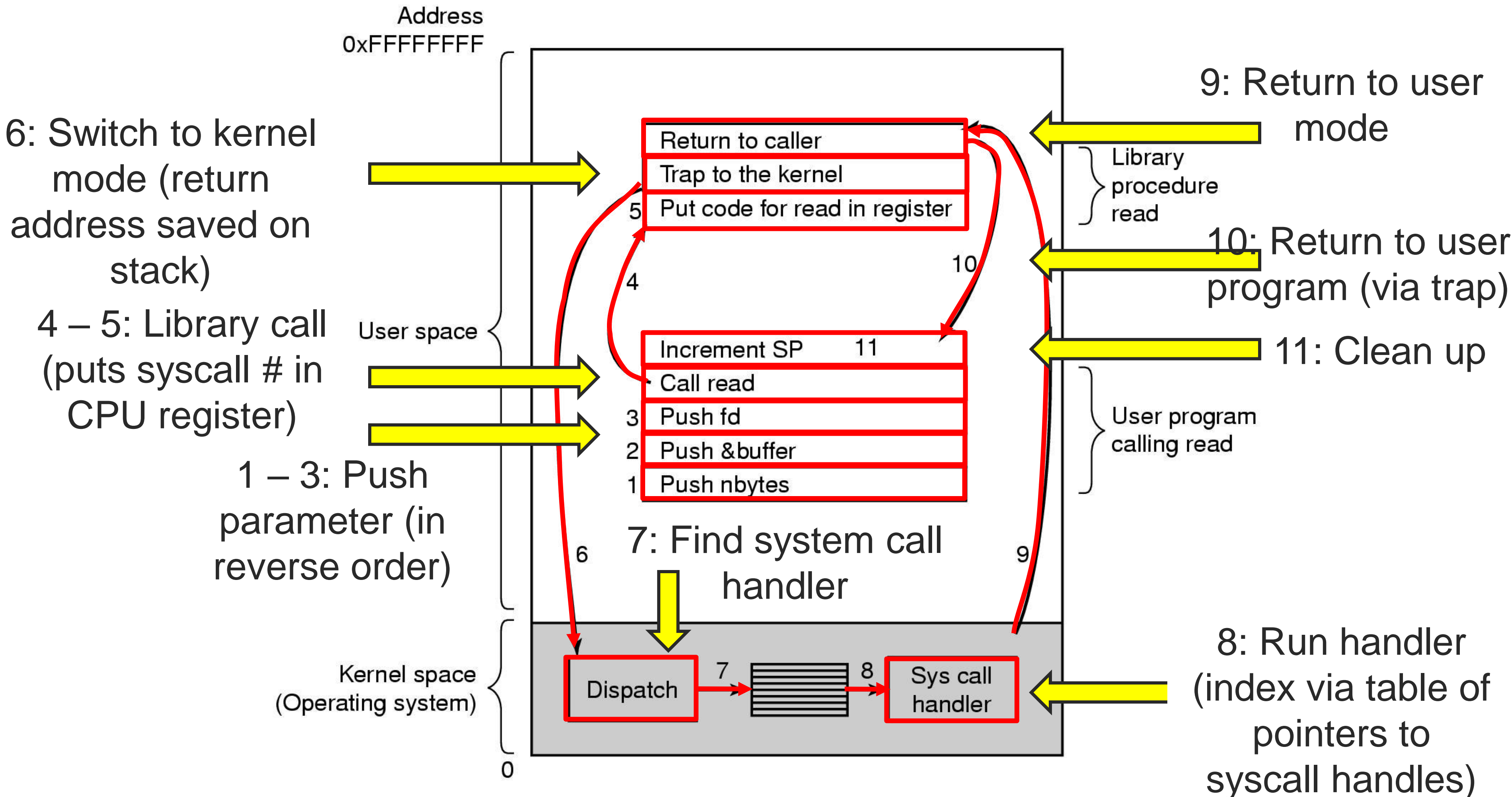
Steps for Making a System Call (Example: read call)



Five categories of system calls

- Process Control
  - `fork()`: Creates a child process
  - `exec()`: Execute a new process image
  - `kill()`: Terminate/signal a process
  - `wait()`: Wait for a process to complete
  - `sbrk()`: Increase process' heap size
  - …

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Five categories of system calls

- File Management
  - `open()`: Opens a file
  - `close()`: Closes a file
  - `read()`: Reads from a file
  - `write()`: Writes to a file
  - `lseek()`: Seek within a file
  - …

Five categories of system calls

- Device Management
  - `mkdir()`: Makes a directory
  - `rmdir()`: Removes an empty directory
  - `link()`: Creates a link to a file/directory
  - `unlink()`: Removes the link
  - `mount()`: Mount a device/file system
  - `unmount()`: Removes the mount
  - …

Five categories of system calls

- Information Management
  - `stat()`: Get status of a file/directory
  - `times()`: Process running times
  - `getrusage()`: Resource usage
  - `clock_gettime()`: Get system time
  - `clock_getres()`: Clock resolution
  - …

Five categories of system calls

- Communication
  - `pipe()`: Communicate b/t two processes
  - `shmget()`: Share memory b/t processes
  - `mmap()`: Maps virtual memory
  - `socket()`: Network socket
  - `connect()`: Connect to a remote server
  - `accept()`: Accept remote connection
  - `send()`: Send network messages
  - …

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System Call Errors…

- When a system call fails, it sets a special global variable: **errno**

Basic Unix Concepts

- Error Model
  - Return value
    - 0 on success
    - -1 on failure for functions returning integer values
    - NULL on failure for functions returning pointers
  - Examples (see **errno.h**)

```
#define EPERM     1       /* Operation not permitted */
#define ENOENT    2       /* No such file or directory */
#define ESRCH     3       /* No such process */
#define EINTR     4       /* Interrupted system call */
#define EIO       5       /* I/O error */
#define ENXIO     6       /* No such device or address */
```

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How do we know what is a system call?

open(2) - Linux man page

Name

open, creat - open and possibly create a file or device

Synopsis

```
#include <sys/types.h>#include <sys/stat.h>#include <fcntl.h>
int open(const char *pathname, int flags);int open(const char *pathname,
int flags, mode_t mode);int creat(const char *pathname, mode_t mode);
```

Description

Given a *pathname* for a file, **open**() returns a file descriptor, a small, nonnegative integer for use in subsequent system calls (*read*(2), *write*(2), *lseek*(2), *fcntl*(2), etc.). The file descriptor returned by a successful call will be the lowest-numbered file descriptor not currently open for the process.

By default, the new file descriptor is set to remain open across an *execve*(2) (i.e., the **FD_CLOEXEC** file descriptor flag described in *fcntl*(2) is initially disabled; the **O_CLOEXEC** flag, described below, can be used to change this default). The file offset is set to the beginning of the file (see *lseek*(2)).

A call to **open**() creates a new *open file description*, an entry in the system-wide table of open files. This entry records the file offset and the file status flags (modifiable via the *fcntl*(2) **F_SETFL** operation). A file descriptor is a reference to one of these entries; this reference is unaffected if *pathname* is subsequently removed or modified to refer to a different file. The new open file description is initially not shared with any other process, but sharing may arise via *fork*(2).

The argument *flags* must include one of the following *access modes*: **O_RDONLY**, **O_WRONLY**, or **O_RDWR**. These request opening the file read-only, write-only, or read/write, respectively.

In addition, zero or more file creation flags and file status flags can be bitwise-*or*'d in *flags*. The *file* creation flags are **O_CREAT**, **O_EXCL**, **O_NOCTTY**, and **O_TRUNC**. The *file* status flags are all of the remaining flags listed below. The distinction between these two groups of flags is that the file status flags can be retrieved and (in some cases) modified using *fcntl*(2). The full list of file creation flags and file status flags is as follows:

die.net

site search

Google Custom Search

Library
- linux docs
- linux man pages
- online dictionary
- page load time

Toys
- world sunlight
- moon phase
- trace explorer

- **2: System Call**
- **3: Library Call**

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Five categories of system calls

- Process Control
- **File Management**
- Device Management
- Information Management
- Communication

File System and I/O Related System Calls

- A file system
  - A means to organize, retrieve, and update data in persistent storage
  - A hierarchical arrangement of directories
  - Bookkeeping information (file metadata)
    - File length, # bytes, modified timestamp, etc
- Unix file system
  - Root file system starts with "/"

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Why does the OS control I/O?

- Safety
  - The computer must try to ensure that if a program has a bug in it, then it doesn't crash or mess up
    - The system
    - Other programs that may be running at the same time or later
- Fairness
  - Make sure other programs have a fair use of device

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File: Open

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
int open (const char* path, int flags [, int mode ]);
```

- Open (and/or create) a file for reading, writing or both
- Returns:
  - Return value ≥ 0 : Success - New file descriptor on success
  - Return value = -1: Error, check value of `errno`
- Parameters:
  - `path`: Path to file you want to use
    - Absolute paths begin with “/”, relative paths do not
  - `flags`: How you would like to use the file
    - `O_RDONLY`: read only, `O_WRONLY`: write only, `O_RDWR`: read and write, `O_CREAT`: create file if it doesn’t exist, `O_EXCL`: prevent creation if it already exists

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File: Open

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
int open (const char* path, int flags [, int mode ]);
```

- Open (and/or create) a file for reading, writing or both
- Returns:
  - Return value ≥ 0 : Success - New file descriptor on success
  - Return value = -1: Error, check value of `errno`
- Parameters:
  - `path`: Path to file you want to use
    - Absolute paths begin with "`/`", relative paths do not
  - `flags`: How you would like to use the file
    - `O_RDONLY`: read only, `O_WRONLY`: write only, `O_RDWR`: read and write, `O_CREAT`: create file if it doesn't exist, `O_EXCL`: prevent creation if it already exists

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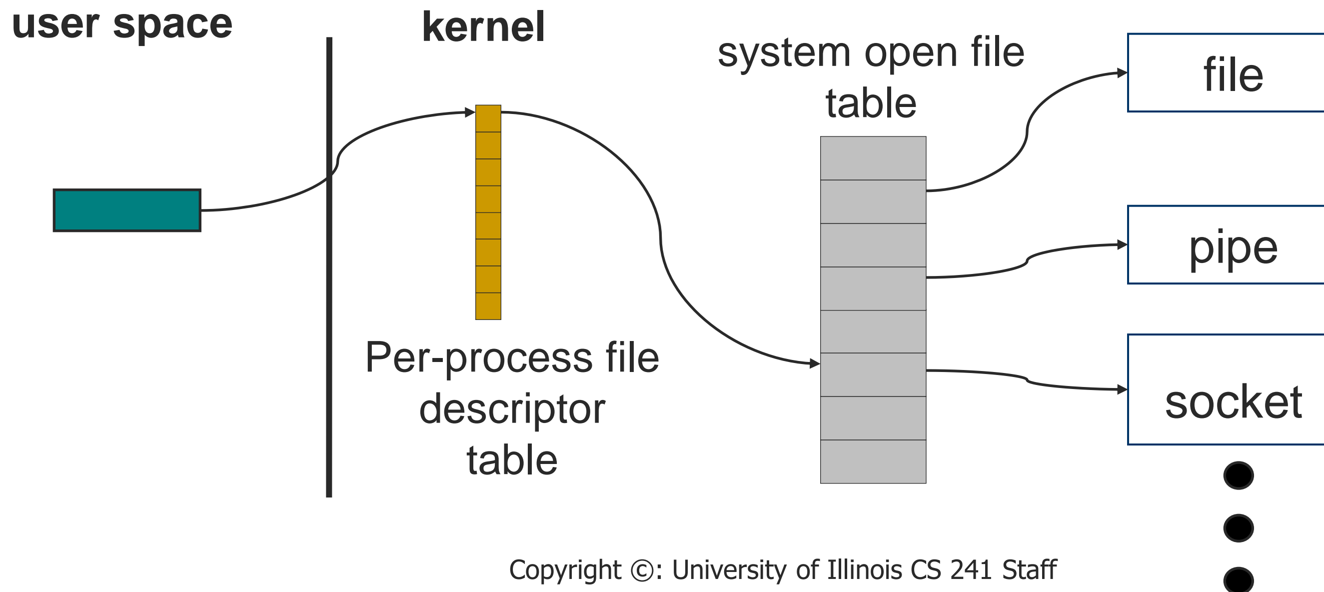
File Descriptors

- Input/Output – I/O
  - Per-process table of I/O channels
  - Table entries describe files, sockets, devices, pipes, etc.
  - Table entry/index into table called “file descriptor”
  - Unifies I/O interface



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System Calls for I/O

- Three file descriptors are always defined:
  - `0 := stdout`
  - `1 := stdin`
  - `2 := stderr`

System Calls for I/O

- Get information about a file

```
int stat(const char* name, struct stat* buf);
```

- Open (and/or create) a file for reading, writing or both

```
int open (const char* name, in flags);
```

- Read data from one buffer to file descriptor

```
size_t read (int fd, void* buf, size_t cnt);
```

- Write data from file descriptor into buffer

```
size_t write (int fd, void* buf, size_t cnt);
```

- Close a file

```
int close(int fd);
```

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System Calls for I/O

- They look like regular procedure calls but are different
  - A system call makes a request to the operating system by trapping into kernel mode
  - A procedure call just jumps to a procedure defined elsewhere in your program
- Some library procedure calls may themselves make a system call
  - e.g., `fopen()` calls `open()`

POSIX I/O vs. C I/O

- `open()`
- `read()`
- `write()`
- `lseek()`
- `close()`

- `fopen()`
- `fread()`
  - `scanf()`
  - `fgetc()`
- `fwrite()`
  - `fprintf()`
- `fseek()`
- `fclose()`

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POSIX I/O vs. C I/O

- `open()`
- `read()`
- `write()`
- `lseek()`
- `close()`

POSIX I/O:
- More powerful functionally
- Only runs on POSIX systems

- `fopen()`
- `fread()`
  - `scanf()`
  - `fgetc()`
- `fwrite()`

C I/O:
- General functionality
- Works on Windows/Linux/etc
- On Linux, calls POSIX I/O

File: Statistics

```
#include <sys/stat.h>
int stat(const char* name, struct stat* buf);
```

- Get information about a file
- Returns:
  - 0 on success
  - -1 on error, sets `errno`
- Parameters:
  - `name`: Path to file you want to use
    - Absolute paths begin with "`/`", relative paths do not
  - `buf`: Statistics structure
    - `off_t st_size`: Size in bytes
    - `time_t st_mtime`: Date of last modification. Seconds since January 1, 1970
- Also

```
int fstat(int filedes, struct stat *buf);
```

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File: Close

```
#include <fcntl.h>
int close(int fd);
```

- Close a file
  - Tells the operating system you are done with a file descriptor
- Return:
  - 0 on success
  - -1 on error, sets `errno`
- Parameters:
  - `fd`: file descriptor

Example (close())

```
#include <fcntl.h>
main(){
    int fd1;

    if(( fd1 = open("foo.txt", O_RDONLY)) < 0){
        perror("c1");
        exit(1);
    }
    if (close(fd1) < 0) {
        perror("c1");
        exit(1);
    }
    printf("closed the fd.\n");
```

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Example (close())

```
#include <fcntl.h>
main(){
    int fd1;

    if(( fd1 = open(“foo.txt", O_RDONLY)) < 0){
        perror("c1");
        exit(1);
    }
    if (close(fd1) < 0) {
        perror("c1");
        exit(1);
    }
    printf("closed the fd.\n");
```

After close, can you still use the file descriptor?

Why do we need to close a file?

File: Read

```
#include <fcntl.h>
size_t read (int fd, void* buf, size_t cnt);
```

- Read data from one buffer to file descriptor
  - Read `size` bytes from the file specified by `fd` into the memory location pointed to by `buf`
- Return: How many bytes were actually read
  - Number of bytes read on success
  - 0 on reaching end of file
  - -1 on error, sets `errno`
  - -1 on signal interrupt, sets `errno` to `EINTR`
- Parameters:
  - `fd`: file descriptor
  - `buf`: buffer to read data from
  - `cnt`: length of buffer

File: Read

```
size_t read (int fd, void* buf, size_t cnt);
```

- Things to be careful about
  - `buf` needs to point to a valid memory location with length not smaller than the specified size
    - Otherwise, what could happen?
  - `fd` should be a valid file descriptor returned from `open()` to perform read operation
    - Otherwise, what could happen?
  - `cnt` is the requested number of bytes read, while the return value is the actual number of bytes read
    - How could this happen?

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Example (**read()**)

```
#include <fcntl.h>
main() {
    char *c;
    int fd, sz;

    c = (char *) malloc(100
              * sizeof(char));
    fd = open("foo.txt",
              O_RDONLY);
    if (fd < 0) {
        perror("r1");
        exit(1);
    }
```

```
sz = read(fd, c, 10);
printf("called
     read(%d, c, 10).
     returned that %d
     bytes  were
     read.\n", fd, sz);
c[sz] = '\0';

printf("Those bytes
     are as follows:
     %s\n", c);
close(fd);
}
```

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File: Write

```
#include <fcntl.h>
size_t write (int fd, void* buf, size_t cnt);
```

- Write data from file descriptor into buffer
  - Writes the bytes stored in `buf` to the file specified by `fd`
- Return: How many bytes were actually written
  - Number of bytes written on success
  - 0 on reaching end of file
  - -1 on error, sets `errno`
  - -1 on signal interrupt, sets `errno` to `EINTR`
- Parameters:
  - `fd`: file descriptor
  - `buf`: buffer to write data to
  - `cnt`: length of buffer

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File: Write

```
size_t write (int fd, void* buf, size_t cnt);
```

- Things to be careful about
  - The file needs to be opened for write operations
  - `buf` needs to be at least as long as specified by `cnt`
    - If not, what will happen?
  - `cnt` is the requested number of bytes to write, while the return value is the actual number of bytes written
    - How could this happen?

Example (write())

```
#include <fcntl.h>
main()
{
    int fd, sz;

    fd = open("out3",
        O_RDWR | O_CREAT |
        O_APPEND, 0644);
    if (fd < 0) {
        perror("r1");
        exit(1);
    }

    sz = write(fd, "cs241\n",
        strlen("cs241\n"));

    printf("called write(%d,
        \"cs360\\n\", %d).
        it returned %d\n",
        fd, strlen("cs360\n"),
        sz);

    close(fd);
}
```

File Pointers

- All open files have a "file pointer" associated with them to record the current position for the next file operation
- On open
  - File pointer points to the beginning of the file
- After reading/write m bytes
  - File pointer moves m bytes forward

File: Seek

```
#include <unistd.h>
off_t lseek(int fd, off_t offset, int whence);
```

- Explicitly set the file offset for the open file
- Return: Where the file pointer is
  - the new offset, in bytes, from the beginning of the file
  - -1 on error, sets **errno**, file pointer remains unchanged
- Parameters:
  - **fd**: file descriptor
  - **offset**: indicates relative or absolute location
  - **whence**: How you would like to use **lseek**
    - **SEEK_SET**, set file pointer to **offset** bytes from the beginning of the file
    - **SEEK_CUR**, set file pointer to **offset** bytes from current location
    - **SEEK_END**, set file pointer to **offset** bytes from the end of the file

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File: Seek Examples

- Random access
  - Jump to any byte in a file
- Move to byte #16

```
newpos = lseek(fd, 16, SEEK_SET);
```

- Move forward 4 bytes

```
newpos = lseek(fd, 4, SEEK_CUR);
```

- Move to 8 bytes from the end

```
newpos = lseek(fd, -8, SEEK_END);
```

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Example (`lseek()`)

```
c = (char *) malloc(100 *
    sizeof(char));
fd = open("foo.txt", O_RDONLY);
if (fd < 0) {
    perror("r1");
    exit(1);
}

sz = read(fd, c, 10);
printf("We have opened in1, and
    called read(%d, c, 10).\n",
    fd);
c[sz] = '\0';
printf("Those bytes are as
    follows: %s\n", c);
```

```
i = lseek(fd, 0, SEEK_CUR);
printf("lseek(%d, 0, SEEK_CUR)
    returns that the current
    offset is %d\n\n", fd, i);

printf("now, we seek to the
    beginning of the file and
    call read(%d, c, 10)\n",
    fd);
lseek(fd, 0, SEEK_SET);
sz = read(fd, c, 10);
c[sz] = '\0';
printf("The read returns the
    following bytes: %s\n", c);
…
```

Standard Input, Standard Output and Standard Error

- Every process in Unix has three predefined file descriptors
  - File descriptor 0 is standard input (**STDIN**)
  - File descriptor 1 is standard output (**STDOUT**)
  - File descriptor 2 is standard error (**STDERR**)
- Read from standard input,
  - `read(0, ...);`
- Write to standard output
  - `write(1, ...);`
- Two additional library functions
  - `printf();`
  - `scanf();`

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Stream Processing - `fgetc()`

```
int fgetc(FILE *stream);
```

- Read the next character from **`stream`**
- Return
  - An *unsigned char* cast to an *int*
  - **`EOF`** on end of file
  - Error

```
int getchar(void);
```

  - Read the next character from **`stdin`**

```
int getc(void);
```

  - Similar to , but implemented as a macro, faster and potentially unsafe

Similar functions for writing:

```
int fputc(int c, FILE *stream);
int putchar(int c);
int putc(int c, FILE *stream);
```

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Stream Processing - `fgets()`

```
char *fgets(char *s, int size, FILE
    *stream);
```

- Read in at most one less than **`size`** characters from **`stream`**
  - Stores characters in buffer pointed to by *s*.
  - Reading stops after an **`EOF`** or a newline.
  - If a newline is read, it is stored into the buffer.
  - A `'\0'` is stored after the last character in the buffer.
- Return
  - **`s`** on success
  - NULL on error or on **`EOF`** and no characters read

Similar:

```
int fputs(const char *s, FILE *stream);
```

Stream Processing

```
char *gets(char *s);
```

- Reads a line from stdin
- NOTE: DO NOT USE
  - Reading a line that overflows the array pointed to by *s* causes undefined results.
  - The use of is `fgets()` recommended

Stream Processing - `fputs()`

```
int fputs(const char *s, FILE *stream);
```

- Write the null-terminated string pointed to by `s` to the stream pointed to by `stream`.
  - The terminating null byte is not written
- Return
  - Non-neg number on success
  - `EOF` on error

```
char *puts(char *s);
```

- Write to `stdout`
  - Appends a newline character

Example: (fgets() - fputs())

```
#include <stdio.h>
int main() {
   FILE * fp = fopen("test.txt", "r");
   char line[100];
   while( fgets(line, sizeof(line), fp) != NULL )
      fputs(line, stdout);
   fclose(fp);
   return 0;
}
```

Stream Processing - `fscanf()`

```
int scanf(const char *format, ... );
```

- Read from the standard input stream **`stdin`**
  - Stores read characters in buffer pointed to by *s*.
- Return
  - Number of successfully matched and assigned input items
  - **`EOF`** on error

```
int fscanf(FILE *stream, const char *fmt, ... );
```

- Read from the named input **`stream`**

```
int sscanf(const char *s, const char *fmt, ... );
```

- Read from the string **`s`**

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Example: (scanf())

- Input: `56789 56a72`

```
#include <stdio.h>
int main() {
   int i;
   float x;
   char name[50];
   scanf("%2d%f %[0123456789]", &i, &x, name);
}
```

What are **i**, **x**, and **name** after the call to **scanf()**?

What will a subsequent call to **getchar()** return?

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Example: `stdin`

```
int x;
char st[31];

/* read first line of input */
printf("Enter an integer: ");
scanf("%d", &x);

/* read second line of input */
printf("Enter a line of text: ");
fgets(st, 31, stdin);
```

What will this code really do?

Example: `stdin`

```
int x;
char st[31];

/* read first line of input */
printf("Enter an integer: ");
scanf("%d", &x);

/* read second line of input */
printf("Enter a line of text: ");
fgets(st, 31, stdin);
```

What will this code really do?

Input is buffered, but `scanf()` did not read all of the first line

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Example: `stdin`

```
int x;
char st[31];
/* read first line */
printf("Enter an
   integer: ");
scanf("%d", &x);
dump_line(stdin);
/* read second line */
printf("Enter a line of
   text: ");
fgets(st, 31, stdin);
```

```
void dump_line( FILE * fp
   ) {
   int ch;
   while((ch = fgetc(fp))
      != EOF &&
      ch != '\n' )
      /* null body */;
   }
```

Read and dump all characters from input buffer until a `'\n'` after `scanf()`